A NICHED PARETO GENETIC ALGORITHM For Multiple Sequence Alignment Optimization

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- Abstract: The alignment of molecular sequences is a recurring task in bioinformatics, but it is not a trivial problem. The size and complexity of the search space involved difficult the task of finding the optimal alignment of a set of sequences. Due to its adaptive capacity in large and complex spaces, Genetic Algorithms emerge as good candidates for this problem. Although they are often used in single objective domains, its use in multidimensional problems allows finding a set of solutions which provide the best possible optimization of the objectives – the Pareto front. Niching methods, such as sharing, distribute these solutions in space, maximizing their diversity along the front. We present a niched Pareto Genetic Algorithm for sequence alignment which we have tested with six BAliBASE alignments, taking conclusions regarding population evolution and quality of the final results. Whereas methods for finding the best alignment are mathematical, not biological, having a set of solutions which facilitate experts' choice, is a possibility to consider.

1 INTRODUCTION

The alignment of protein, DNA and RNA sequences is a very frequent task in bioinformatics. Multiple sequence alignment is an optimization problem which consists on finding the best alignment from large complex search spaces (Horng et al., 2005). Its main goal is to help in the comparison of sequence structure relationship, by identifying sequences' similarities and differences (Pal et al., 2006).

Genetic Algorithms (GAs) are search algorithms based on the principals of natural evolution and genetics (Goldberg, 1989). They are able to take advantage of gathering information about an initially unknown search space, in order to bias subsequent search into useful subspaces. This quality makes them suitable for problems with large, complex, and poorly understood search spaces (De Jong, 1988), such as multiple sequence alignment. Although GAs are often used in single objective problems, they can also be used in multiobjective problems, on which the GA is used to find all possible tradeoffs among the multiple conflicting objectives (Horn et al., 1994). The resulting non-dominated solutions lie on the Pareto optimal frontier, meaning that there are no other solutions superior in all objectives.

Niching methods, such as sharing, helps in maintaining the diversity of certain properties within the population, preventing the convergence to a single point in the Pareto front and allowing parallel convergence into multiple good solutions (Shir and Back, 2006).

In our prior investigation we have developed AlineaGA, a genetic algorithm which performs multiple sequence alignment. In our first approach, we tested AlineaGA with a single objective fitness function - the sum-of-pairs (Silva et al., 2008). Later, we tested the weighted sum of the sum-of-pairs value with the number of fully identical columns to perform alignment evaluation (Silva et al., 2009). Now, we present a multiobjective strategy which tries to maximize both the sum-of-pairs and the number of fully identical columns by means of a niching mechanism named equivalence class sharing (Horn et al., 1994). Our objective is to evaluate the quality of the found solutions using this approach. For this matter, we have tested AlineaGA with six BAliBASE (Thompson et al., 1999) alignments.

José Mateus da Silva F., Manuel Sánchez Pérez J., Antonio Gómez Pulido J. and A. Vega Rodríguez M. (2010). A NICHED PARETO GENETIC ALGORITHM - For Multiple Sequence Alignment Optimization. In Proceedings of the 2nd International Conference on Agents and Artificial Intelligence - Artificial Intelligence, pages 323-329 DOI: 10.5220/0002729303230230 This paper is organized as follows. In the next Section we introduce concepts underlying our research. In Section 3, we present a brief explanation regarding AlineaGA methods. Section 4 presents AlineaGA's niched Pareto approach. The experiments performed in order to observe the impact of these strategy are discussed in Section 5. Finally, the concluding Section presents final considerations and topics for future work.

2 BACKGROUND

Although it may not be obvious, multiple sequence alignments are present in most of the computational methods used in molecular biology. They are used in different areas such as functional genomics, structure modelling, mutagenesis experiments, evolutionary studies and drug design.

There are several approaches to the sequence alignment. The two most important ones are based on progressive and iterative methods.

When progressive methods are used, the alignment is gradually built up by aligning the two most similar sequences first, and adding the less similar ones one after another. This fast and simple method has a critical problem: if a mistake is made at an intermediate step, it cannot be corrected later by adding the remaining sequences. Also, it does not provide a metric which allows the comparison of two different alignments of the same set of sequences, or which can be used to say that the best possible alignment, for a set of parameters, have been found (Notredame and Higgins, 1996).

Iterative methods try to optimize a scoring function which reflects the biological events which took place in the evolution of the sequences. Optimizing this score leads to a correct alignment (Lassmann and Sonnhammer, 2002). One example of iterative methods are GAs, other examples may be found in our prior review (Silva et al., 2007).

2.1 Alignment

An alignment is an arrangement of two or more sequences in a way which reveals where the sequences are similar, and where they differ. An optimal alignment exhibits the most correspondences and the fewest differences, even if it will not be biologically meaningful (Pal et al., 2006). Figure 1 shows an example of an alignment of four hypothetical protein sequences.

-TISCTGNIGAG-NHVKWYQQLPG
-RLSCSSIFSSYAMYWVRQAPG
L-LTCTVSFDDYYSTWVRQPPG

Figure 1: Example of a multiple sequence alignment.

Sequences may have different lengths and each one is represented in a different line. Columns with the same characters, presented in bold, show that in that specific position, no mutation occurs among the sequences. On the other hand, columns which present different characters show that mutation events have taken place. The characters used to represent the elements of the molecular sequences are often referred as residues.

Gaps can be introduced in the sequences, allowing the alignment to be extended into regions where its sequences may have lost or gained residues. These gaps are usually represented by the symbol "–".

2.2 Genetic Algorithms

GAs, are a class of evolutionary algorithms introduced by Holland (Holland, 1975). Its search methods model some natural facts: genetic inheritance and Darwinian strife for survival (Michalewicz, 1996).

In GAs, the adaptation is done by keeping a population of structures from which new structures are produced through genetic operators, such as crossover and mutation(De Jong, 1988).

In crossover, characteristics of two randomly chosen individuals (parents), are combined to form two similar offspring by swapping corresponding segments of parents. Mutation randomly alters some values within the individual by a arbitrary change (Anbarasu et al., 2000). Each structure of the population has a fitness score, which is used to choose which structures will be used to form new ones (De Jong, 1988).

The ability to gather information about a search space, initially unknown, to direct the search for useful subspaces, is a distinguishing characteristic of GAs. This ability makes them suitable for solving problems with large, complex and unknown search spaces (De Jong, 1988).

2.3 Fitness Sharing

Fitness sharing (Goldberg and Richardson, 1987) is a mechanism for maintaining population diversity. It distributes the population over different peaks in the search space by reducing the fitness of highly similar solutions. Equation 1 presents the shared fitness of an individual *i*, where f_i is the individual raw fitness and m_i is the nich count, representing how crowded is the neighborhood of individual *i*.

$$f_i^{share} = \frac{f_i}{m_i} \tag{1}$$

The nich count is computed by adding a sharing function over all members of the population as follows:

$$m_i = \sum_{j=1}^n Sh\left(d_{i,j}\right) \tag{2}$$

Where $Sh(d_{i,j})$ represents the sharing function, presented in Equation 3, and $d_{i,j}$ is the distance between the *i* and *j* individuals, which can be based on either phenotype or genotype similarity.

$$Sh(d_{i,j}) = \begin{cases} 1 - \frac{d_{i,j}}{\sigma_{share}} & \text{if } d \le \sigma_{share} \\ 0 & \text{if } d \ge \sigma_{share} \end{cases}$$
(3)

The niche radius is given by σ_{share} . Solutions within this radius are in the same neighborhood, reducing each other's fitness.

3 AlineaGA METHODS

In AlineaGA, the initial population is randomly generated, and then the individuals are selected, combined and mutated in order to produce new solutions through the course of a defined number of generations. This section presents a brief explanation regarding AlineaGA's representation, evaluation, crossover and mutation.

3.1 Representation

We use a non-codified representation of the individuals. Real multiple sequence alignments, as the one presented in Figure 1, are used as data structures for each individual. Chromosomes are represented by arrays of characters on which each line corresponds to a sequence in the alignment, and each column represents a residue at a specific position.

3.2 Evaluation

To perform the evaluation of each solution, two attributes are used: the sum-of-pairs and the identity of the alignment. The sum-of-pairs function, presented in Equation 4, is assessed by scoring all of the pairwise comparisons between each residue in each column of an alignment and adding the scores together (Wang and Lefkowitz, 2005).

$$Sum - of - Pairs = \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} ScoringMatrix(l_i, l_j)$$
(4)

For this purpose, a scoring matrix which determines the cost of substituting a residue for another is used, as well as a gap penalty value to determine the cost of aligning a residue with a gap. We use the PAM 350 (Dayhoff et al., 1978) scoring matrix with a gap penalty of -10 (Silva et al., 2008).

The identity of the alignment is simply the number of fully identical columns in the alignment.

3.3 Crossover

AlineaGA uses one of the two crossover operators, randomly selected within each generation. The One Point crossover derives from Goldberg's standard one point crossover operator (Goldberg, 1989) with an extension that treats the existing gaps in each sequence. On RecombineMatchedCol (Chellapilla and Fogel, 1999), the fully identical columns of the first parent which do not appear in the second one are identified, and then, one of these fully aligned columns is randomly selected and is generated in the second alignment, originating the offspring.

3.4 Mutation

Each mutation operator is randomly selected from a pool of six operators and it is applied to an individual according with the defined mutation probability. Whenever the mutated solution is worst than the original one, a new mutation must be applied to the mutated individual. This process is repeated until the fitness improves or during a specific number of attempts. We opted for the maximum of 2 tries. This strategy allows a good tradeoff between speed and robustness, without transforming completely the solutions in a single generation.

The Gap Insertion operator extends the alignments by inserting gaps into the sequences in a random fashion, such as in GenAlignRefine (Wang and Lefkowitz, 2005) gap insertion operator.

Shifting gaps is another way to introduce new alignment configurations. In the Gap Shifting mutation operator, a gap is randomly chosen in an alignment and it is moved to another position in the same sequence (Notredame et al., 1997).

The Merge Space operator merges together two or three spaces of a sequence (Horng et al., 2000). It randomly selects two or three consecutive gaps of a sequence, adjacent or not adjacent, and then merges these gaps together. After that, they are shifted to a randomly chosen position in the same sequence. The Smart Merge Space is similar to the Merge Space operator, but it only applies the mutation if the fitness of the mutated solution is greater than the fitness of the original one (Silva et al., 2009).

The Smart Gap Insertion is a variation of the Gap Insertion operator which only produces the mutation when the fitness of the mutated alignment is greater than the fitness of the original one (Silva et al., 2008). The insertion of additional gaps is determined by a direction probability which reflects the success of inserting gaps at the beginning or at the end of the alignment. If the operator does not improve the alignment at the first attempt, it chooses a new random position of insertion and repeats the whole process. The defined number of maximum attempts is set to 3, but it can be customized according to user's needs.

The Smart Gap Shifting, tries to move the gaps of an alignment until its fitness improves (Silva et al., 2008). As in the Smart Gap Insertion operator, the shift direction is determined by a direction probability which is updated when better alignments are found. Likewise, the mutation occurs only if the fitness of the generated alignment is greater than the original one.

The use of crossover and mutation operators can produce columns completely formed by gaps in the alignment. To remove these gap columns we use the Gap Column Remover (Silva et al., 2008), which is not conditioned by the mutation probability and it is applied at the end of each generation.

4 NICHED PARETO GA

The Niched Pareto GA is characterized by its selection mechanism. In previous works (Silva et al., 2008, Silva et al., 2009), we use tournament selection to choose the solutions of the current generation that will prevail for the next one. However, throughout the generations, this technique tends to lead the population to a single point in the search space. To maintain multiple Pareto optimal solutions and avoid convergence, we use Pareto domination tournaments and equivalence class sharing (Horn et al., 1994), which we now present.

4.1 Pareto Domination Tournaments

In a normal binary tournament, two randomly selected individuals compete for domination. If one dominates the other, it wins. However, this condition does not produce a sufficient domination pressure. Pareto domination tournaments (Horn et al., 1994)

use a sampling scheme which offers control over the domination pressure. In this method, two candidate solutions are randomly chosen from the population for selection purposes. Also, a comparison set is formed by randomly choosing individuals from the population. Then, each candidate solution is compared with every individual in the comparison set. The candidate which dominates all the individuals in the comparison set is selected for reproduction. If both candidates dominate or are dominated by the comparison set, then sharing is used to select the winner, as section 4.2 explains.

Adjusting the size of the comparison set allows the control of the domination pressure. High values for this parameter tend to increase the pressure towards a small portion of the front. On the other hand, small comparison sets result in many dominated solutions. Typically, a comparison set with size of 10% of the population, yields a tight and complete distribution over the front (Horn et al., 1994).

4.2 Equivalence Class Sharing

To avoid genetic drift, whenever the candidate solutions are both dominated or both non-dominated by the comparison set, the winner is selected by equivalence class sharing (Horn et al., 1994).

This particular method of sharing does not degrade the fitness of the individuals. Instead, it assumes that candidates, mutually dominated or non-dominated, are equally fit. Therefore, in order to maintain diversity along the Pareto front, this method computes the nich count of both candidates and selects the one which has the smallest number of individuals on its neighbourhood.

4.2.1 Distance Metric

The distance metric may be based on either phenotype or genotype similarity. In our particular case, the genotype and phenotype representation are the same. As we are trying to maximize two different objectives represented in a 2 dimensional space, we opt for using the Euclidean distance as a similarity measure.

4.2.2 Niche Radius σ_{share}

Defining the radius which determines each nich range is not a trivial mater. Such as (Shir and Back, 2006), we determine the σ_{share} value according with Equation 5.

$$\sigma_{share} = \frac{r}{\sqrt[n]{q}} \tag{5}$$

Datasat	BAliBASE		Number of	AlineaGA			
Dataset	SOP	ID	Peaks	Avg. Best SOP	Avg. Best ID	Best SOP	Best ID
laho			81	1974,83	10,90	2155	13
	2015	12	49	1974,60	11,03	2141	13
			4	1960,03	10,93	2112	13
1 fmb			36	1817,03	24,97	1864	27
	1706	25	100	1811,07	24,93	1860	27
			4	1807	25,40	1864	27
1plc			4	2356	17,33	2590	20
	2403	18	25	2353,87	17,60	2589	20
			100	2340,60	17,10	2576	20
1hpi			4	1135,43	12,17	1198	14
	1208	10	81	1128,30	12,37	1198	14
			36	1120,17	12,64	1201	15
1pfc			16	2442,97	14,23	2519	15
	2216	13	4	2435,90	14,33	2536	17
			49	2425,17	14,17	2533	16
lycc			36	883,93	6,9	1091	10
	963	11	9	864,03	7,2	1093	10
			64	859,47	6,7	1045	11

Table 1: Results for the AlineaGA Niched Pareto test configurations.

SOP, sum-of-pairs; ID, identity; Avg., Average. Avg. Best SOP and Avg. Best ID were obtained by averaging the results of 30 runs.

The existing theory for setting this value, assumes that the solution set has a previously known finite number of peaks q (Shir and Back, 2006).

By knowing the upper and lower bounds of each objective, r is defined as follows:

$$r = \frac{1}{2} \sqrt{\sum_{k=1}^{n} (x_{k,msx} - x_{k,\min})^2}$$
(6)

Where n defines the number of objectives, which in our particular case, is 2.

The lower and upper bounds of each dimension are computed on every generation, presenting different values as population evolves. However, in multiple sequence alignment, there is no practical way of knowing the maximum number of peaks beforehand. Therefore, we opt to test several values for this parameter, as next section describes.

5 TESTING AND RESULTS

Our goal is to find the best possible solutions which maximize the sum-of-pairs and the identity of each alignment. We test the sharing function with different σ_{share} values, which are obtained by computing the nich radius for various peak values.

In our tests, we use six datasets from the Reference 1 alignments of BAliBASE (Thompson et al., 1999). Three of these datasets (1aho, 1fmb, 1plc,) have more than 35% of identity among its sequences; and the rest (1hpi, 1pfc, 1ycc) present 20% to 40% of identity. We have measured the sum-of-pairs score and the identity of each one of these datasets. Later we use these reference results to evaluate the different test configurations.

5.1 Test Configurations

Although we have tested all our datasets for 4, 9, 16, 25, 36, 49, 64, 81 and 100 peaks, we only present the results for the 3 configurations which obtained the best results on each dataset. Also, we have started by executing the algorithm during 10000 generations with a mutation probability of 0.05, but we have realized that an equivalent final solution set could be achieved in 2000 generations in less time, by increasing the mutation probability to 0.4. Therefore, we have opted for this latter setting. The remaining parameters are the same in all configurations: the population size is 100, the crossover probability is 0.8 and the number of inserted gaps by the Gap Insertion and Smart Gap Insertion operators is 10. Finally, the size of the comparison set for the Pareto domination tournaments is set to 10.

5.2 Results

Next we present the results of tests performed. All the results were obtained by averaging the sum-of-pairs and the identity scores, from 30 runs of AlineaGA, for each configuration/dataset.

5.2.1 Performance

Table 1 summarizes the performance of the top 3 configurations for each test dataset. The "SOP" of BAliBASE alignment column, presents the sum-of-pairs score for the different datasets. This value was computed using the PAM 350 scoring matrix and a gap penalty of -10. The "ID" of BAliBASE shows the number of fully aligned

columns on each BAliBASE's alignment. Columns "Avg. Best SOP" and "Avg. Best ID", show the average sum-of-pairs and the average identity scores obtained in 30 runs of AlineaGA. The best values found for the sum-of-pairs and identity scores are presented in columns "Best SOP" and "Best ID".

As the results state, it is not possible to establish a direct relation between the number of peaks and the percentage of identity of the alignments. This parameter is directly related with each particular alignment and can not be determined in such generic way. Comparing with the BAliBASE alignments, and with the exception of 1hpi dataset, it is possible to find equal or higher values for both objectives simultaneously in our results. However, the average sum-of-pairs and average identity of the 30 executions of each test are superior only in 1fmb and 1pfc datasets.

5.2.2 Population's Evolution

Figures 2 to 7, present the population's fitness evolution for the best configurations on each dataset.

These values were obtained by averaging each solution's sum-of-pairs and identity scores from the 30 runs of the program. Each figure shows the representation of the population throughout the generations in 4 particular moments: generations 500, 1000, 1500 and 2000 - the final solution set.

We can observe that high values for one of the objectives, will necessarily lower other objective' score. Also, after 2000 generations, we can see that the majority of the population is tightly distributed along the front. Nevertheless, there are a few dominated solutions. These solutions result of crossover and mutation, but generally, they are not held. Dataset 1pfc, shown in Figure 6, presented the most atypical evolution, with the resulting front solutions distributed in a small space on which could have featured some individuals with higher identity values present in generation 1500.

6 CONCLUSIONS

By using a multiobjective approach in this domain, we try to offer a solution to a very significant limitation of multiple sequence alignment: its mathematical approach. As stated before, the best alignment is the one which presents the most correspondences and the fewest differences, but which may or may not be biologically meaningful knowledge is needed to validate the results of an alignment tool. By presenting a set of solutions instead of a single one, it is possible for a biologist to observe several hypotheses and so choose the one which is closer to the biological reality.



Figure 2: Population average fitness for 1aho, 81 peaks.



Figure 3: Population average fitness for 1fmb, 36 peaks.



Generation: 500
Generation: 1000
Generation: 1500
Generation: 2000

Figure 4: Population average fitness for 1plc, 4 peaks.



Figure 5: Population average fitness for 1hpi, 4 peaks.



Generation: 500
Generation: 1000
Generation: 1500
Generation: 2000

Figure 6: Population average fitness for 1pfc, 16 peaks.



Figure 7: Population average fitness for 1ycc, 36 peaks.

The main drawback of this method, as it is implemented, is its dependence of previously knowing the expected number of peaks in the search space. This problem may be overcome by trying to identify the number of peaks in the population dynamically, or by using a different approach when computing the nich radius, σ_{share} .

Alternative objectives, such as minimizing the number of gaps, may be used instead of maximizing the identity. However, this kind of approach may have poor results when several gaps are needed to maximize the similarity among the sequences. A possible solution is to increase the complexity of the problem by optimizing three objectives: maximize identity and sum-of-pairs scores, and minimize the number of gaps in the alignment.

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