

# Quality of Experience Evaluation for Data Services over Cellular Networks

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**Abstract:** This paper presents an end-to-end service performance evaluation method that is able to estimate both the Quality of Service (QoS) and Quality of Experience (QoE) associated to different data services over cellular networks. A set of performance indicators are evaluated at each layer of the terminal's protocol stack following a bottom-up process from the physical layer up to the application layer. Then, specific utility functions for each data service are used to map QoS into QoE in terms of Mean Opinion Score (MOS). Three different data services (web browsing, video YouTube and Skype-based voice over IP) have been evaluated in this paper under different network and terminal configurations. Performance results show that the MOS associated to a particular data service is largely affected by the radio level performance (error rate, throughput and delay), so proper protocols' configuration is a key issue to maximize the QoE.

## 1 INTRODUCTION

Next generation mobile communication systems will support diverse types of services across different types of wired/wireless access technologies. The end-to-end Quality of Service (QoS) provision in such a heterogeneous scenario is one of the main topics in networks research nowadays.

The estimation of the service performance and Quality of Experience (QoE) perceived by the user plays a very important role in wireless networks, as it can be a very valuable input for network design, dimensioning, planning, optimization, configuration or upgrade. However, the assessment of the QoE requires analyzing the performance of the whole network (from user equipment to application server or remote user equipment), which includes the following aspects: individual performance figures for each network element, interfaces and interactions between them, protocols behavior, and how the end-user perception is affected by network-related degradations. In addition, the end-user reacts in a different manner to degradations for different services, e.g., the end-user perception is highly affected by the end-to-end delay in conversational services whereas it has a lower impact on background services such as files transfer.

A common issue from network operators' viewpoint is the process of assessing and managing the QoS of their new services as well as evaluating the quality experienced by the end user. Traditionally, network metrics like accessibility, retainability and quality were sufficient to evaluate the user experience for voice services. However, for data services, the correlation between network performance indicators and application performance indicators is not so straightforward due to the following reasons: firstly, data systems have several protocol layers; and secondly, radio data bearers are typically shared among different services and applications. In these conditions, data service performance assessment is usually performed through active terminal monitoring over real networks. Obviously, if the operator wants to collect statistics on a reasonable number of terminals, applications and locations, this process is very expensive and time consuming.

In some cases, the service and/or specific network to be evaluated are not available or cannot be tested and configured easily. In other cases, the service or specific network is available, but it is needed to estimate their performance under specific configurations, scenarios or network conditions that cannot be easily reproduced. Additionally, the

performance subjectively perceived by the end user, i.e. in terms of Mean Opinion Score (MOS), cannot be directly obtained based on network performance metrics. This is due to the behavior of different protocols and mechanisms along the network elements and their protocol stack, as well as the complex translation of QoS metrics to QoE perceived by the user, which is very service dependent. Typically, the QoE has been measured by performing subjective tests to a wide set of users in order to know their satisfaction degree through a MOS indicator, which can range from 1 (bad) to 5 (excellent); this type of methods is obviously costly and time-consuming for both the subscribers and the operator.

A particular application of this type of solution (i.e. mapping network into application performance indicators and MOS) for determining the quality of experience for on-line gaming traffic is described by Gustafsson et al. (2010), with the peculiarity of using a modeling unit to map game and transport level measurements into MOS values. However, this work is only applicable to online-gaming services as the model is just based on game and transport parameters, with no possibility of using performance indicators from lower layers (e.g. radio protocols in a 3G cellular network like MAC, RLC or PDCP). This means that the network element in charge of monitoring the game and transport performance parameters must have access to the application and transport levels.

Other works have focused on the design and/or configuration of lower layers to optimize upper layers' performance (Luo et al., 2000); (De May et al., 2005); (Lassila and Kuusela, 2008), propose new radio resource management techniques which are adaptive to the QoS or QoE (Piamrat et al., 2010) or focus on particular algorithms to enhance objective quality evaluation of a specific service like voice (Lee et al., 2009). However, none of the previous works provides a method to easily evaluate the application layer performance or the QoE for different packet data services over any network configuration.

In this paper we present an end-to-end evaluation method that is able to assess the QoS and QoE for different multimedia applications like video, voice over IP (VoIP) and web-browsing. The proposed framework provides a set of performance indicators like throughput, delay, and loss rate at different points of the whole protocol stack. This approach may be used for different purposes like e.g. estimation of the QoE for new data services over a specific wireless network (this process can help on

the design and optimization of new services in order to improve the QoE). In addition, it provides a good understanding of how the application performance is affected by the end-to-end network behavior and a way to find the most critical layer in the protocol stack without the need of real networks monitoring.

The remainder of this paper is organized as follows. The proposed methodology to evaluate the QoS and QoE associated to any data service is described in section 2. Section 3 presents a set of performance results for three different data services (web browsing, video YouTube and Skype-based VoIP) under different network and terminal configurations. Finally, some concluding comments are given in section 4.

## 2 QUALITY EVALUATION METHODOLOGY

Packet data services performance and end-user's experience can be characterized considering the cumulative performance degradation along the different network elements and protocol stacks plus the effect of the subjectivity and the perception of the end-user. Generally, such performance is assessed through indicators that are very service dependant, such as response time in web browsing or average throughput when downloading a file.

We propose a new methodology for estimating the QoS and QoE perceived by the user for different packet data services over wireless networks. The proposed methodology is based on network and protocol models, service-related parameters and utility functions that map QoS objective metrics into the subjective experienced quality as perceived by the end-user. Such approach allows easily predicting the performance of different services under specific wireless environments (GPRS, UMTS, LTE, etc.) without the need of running, capturing and analyzing the traffic generated from a real scenario. However, the models can be optionally fed from radio and network performance indicators obtained from different sources: a) Network Operation and Support Subsystem (OSS), b) real measurements obtained at the network or the terminal side (if available), or c) simulation results.

The method herein described is based on theoretical models, including the impact of the:

- network elements along the end-to-end path (e.g. user equipment, base station, gateways, server, etc.), protocols and interfaces;
- particular service under analysis, including

aspects like content sizes (e.g. in a web browsing service: web page text size, number of embedded objects, object sizes), protocols and signaling, and specific application performance indicators (APIs);

- end-user perception, which includes how the subjective experienced quality is affected by APIs (e.g. initial delay, total response time) and perception (e.g. resolution).

### 2.1 Quality of Service Evaluation

The model herein proposed offers a methodology of analysis and evaluation of the QoS based on layers. Each layer is modeled and evaluated based on a set of performance indicators. The goal of the proposed methodology is to provide a performance indicator for different services based on the network performance indicators as well as on their own service parameters. The simulations focus on a hypothetical user which experiences a given set of MAC-level radio access performance parameters, from which the model is able to derive application-level QoS indicators after modeling all the intermediate layers. For the topic at hand, a general LTE network architecture has been considered (see Figure 1). In the proposed methodology each layer is modeled and evaluated based on a set of performance indicators. Note that depending on the particular layer, the scope of each performance indicator may include the end-to-end network (for application, transport and network layers) or just the radio interface (for radio specific layers).

The modeling methodology follows a bottom-up approach, from the physical up to the application layer, taking into account the effects with a higher impact on the overall QoS. Therefore, layer  $i$  ( $L_i$ ) provides a set of performance indicators to the layer above ( $i + 1$ ), and successively, up to the application layer (see Figure 1).

Without loss of generality, the following performance indicators are considered as the most relevant and are provided at each protocol layer:

- Transmission rate ( $R_{L_i}$ ): defines the amount of data correctly transferred at layer  $i$  in a given time (in bits per second). The transmission rate will vary at each layer due to different factors, such as protocol headers, packet loss rate, number of retransmissions, etc.
- Delay ( $D_{L_i}$ ): represents the average time (in seconds) that a data unit (at layer  $i$ ) takes to be transported from peer to peer. The delay is a very important indicator for real-time services and also for those services that use reliable and congestion-aware protocols like Transmission Control Protocol (TCP).
- Loss rate ( $P_{L_i}$ ): represents the loss rate of data units at layer  $i$ . This loss rate may be due to errors at the radio interface or data losses at network queues. In general, the impact of data losses can be minimized by applying correction techniques and/or retransmissions at different levels. However, a high loss rate will produce a large number of retransmissions, which reduces the effective information transmission rate.

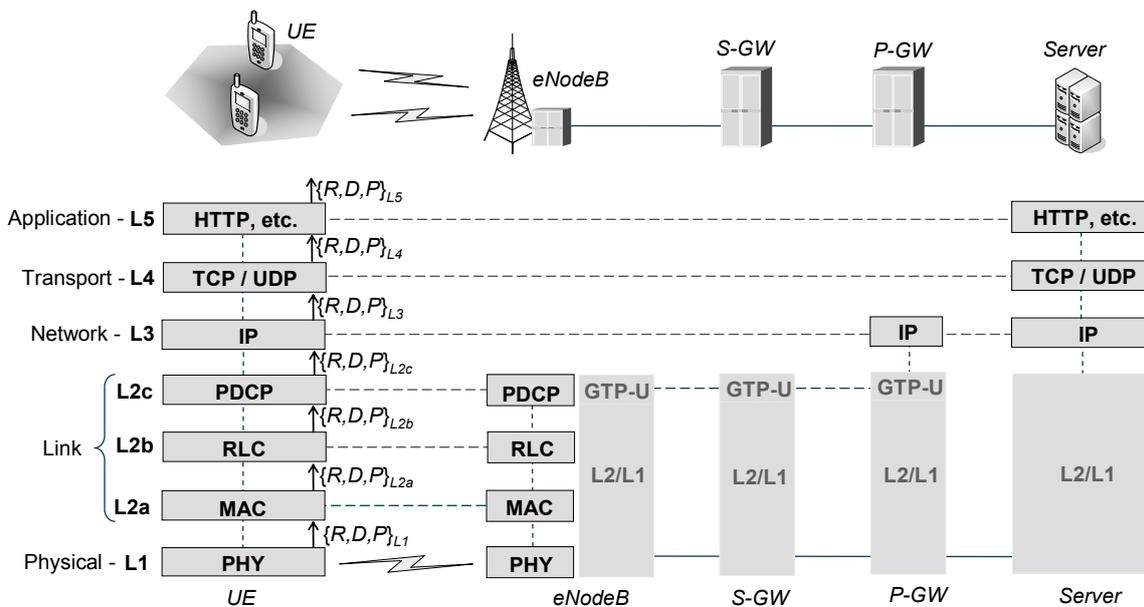


Figure 1: Scenario and protocol stack under analysis.

The final model is composed of a set of deterministic equations starting from the RLC level up to the application level, where performance indicators at layer  $i$  are analytically derived as a function of the performance indicators at layer  $i-1$ . The specific equations that model each layer along the protocol stack is out of the scope of this paper, although further details can be found in a previous work from one of the authors (Gómez et al., 2010b), and a brief summary of the main aspects affecting the QoS at each layer is described in Table 1.

Table 1: Summary of the main aspects affecting the QoS.

Layer		Impact on QoS
Application		The application layer mainly includes the signaling or request/response messages associated to each particular data service.
Transport		The most problematic transport protocol over wireless networks is TCP. Congestion and flow control mechanisms included by TCP have a very negative impact on the throughput and delay, especially for high Round Trip Times (RTTs) and loss rate.
Network		The main aspects affecting the QoS are related to the network RTT and packet loss rate along the end-to-end path.
Link	PDCP	The main impact of PDCP layer on the QoS is due to the use of robust header compression (ROHC), whose gain will be higher as the packet size decreases.
	RLC	It is responsible for the segmentation and reassembly of upper layer data units and, additionally, for performing optional selective retransmissions. Thus, the error rate can be lowered by means of retransmissions at the expense of decreasing the throughput and increasing the average delay and jitter.
	MAC	The MAC layer at the access node allocates channels to users on a subframe basis; that is, for each new subframe, the system assigns available physical channels to users according to a scheduling policy.
Physical (PHY)		Defines the physical channels structure through which the information will be transported.

In this paper we use PHY/MAC link level simulations associated to a LTE network to obtain MAC level performance results under specific configuration (as described in section 3). Such results at the MAC layer are then mapped into performance results at each layer above up to the application layer. Anyhow, MAC layer results from simulations could be replaced by network operator statistics generally available at their OSS database.

## 2.2 Quality of Experience Evaluation

The final goal of this end-to-end model is to evaluate the application level QoS, which will be later mapped into QoE (in terms of MOS value), as

shown in Figure 2. This last process is proposed to be performed by means of utility functions associated to each particular service. The goal of the utility functions is to map objective measurements (in terms of QoS) into subjective metrics (in terms of QoE perceived by the user).

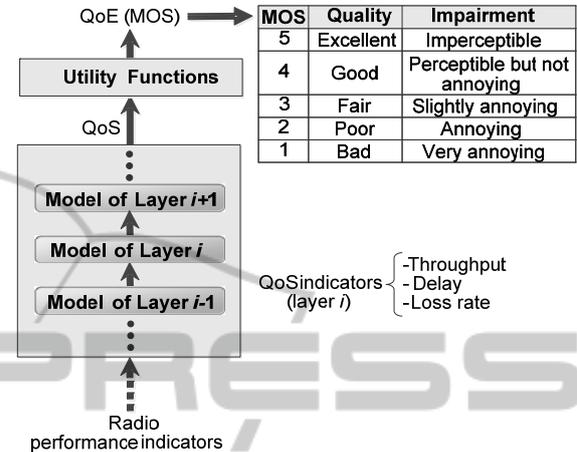


Figure 2: Bottom-up approach to evaluate the QoE.

This mapping process shall consider the specific characteristics of each data service:

- **Web browsing:** the most important objective parameter to estimate the MOS in a web browsing session is the service response time  $D_{L5}$ . The utility function that estimates the MOS as a function of  $D_{L5}$  (in seconds) is given by (Ameigerias, 2010):

$$MOS = 5 - \frac{578}{1 + \left( 11.77 + \frac{22.61}{D_{L5}} \right)^2} \quad (1)$$

- **Video YouTube:** among the various works devoted to estimate the MOS for video services (Mok et al., 2011); (Porter and Peng, 2010); (Ketykó et al., 2010), the analysis presented by Mok et al. (2011) provides a utility function for HTTP video streaming as a function of three application performance metrics: initial buffering time  $T_{init}$  (time elapsed until certain buffer occupancy threshold has been reached so the playback can start, measured in seconds), mean rebuffering time  $T_{rebuf}$  (average duration of a rebuffering event, measured in seconds) and re-buffering frequency  $f_{rebuf}$  (frequency of interruption events during the playback, measured in seconds<sup>-1</sup>). The final MOS expression is given by:

$$MOS = 4.23 - 0.0672T_{init} - 0.742f_{rebuf} - 0.106T_{rebuf} \quad (2)$$

Note that these application layer metrics ( $T_{init}$ ,  $T_{rebuf}$ ,  $f_{rebuf}$ ) can be estimated (at the receiver) from

performance indicators at lower layers (like the TCP throughput) provided by the end-to-end model (described in section 2.1) as well as other configuration parameters like video coding rate or buffer size at the receiver.

▪ **Skype-based VoIP:** in this case the MOS formula just maps the result given by an intermediate model into normalized MOS values. This intermediate model, known as the E model, is specified in ITU-T G.107 (2009) and it provides a numerical estimation  $R \in [0, 100]$  of the voice quality from a set of network impairment factors related with the Signal to Noise Ratio (SNR) of the transmission channel, delay, distortions introduced by the coding/decoding algorithms, packet losses, etc. Cole and Rosenbluth (2001) provides a simplification of the E-model, particularizing it for VoIP communications, where the voice quality  $R$  is:  $R = 94.2 - 0.024 \cdot d - 0.11 \cdot (d - 177.3) \cdot H(d - 177.3) - I_{e-eff} + A$  being  $d$  the end-to-end delay in milliseconds,  $I_{e-eff}$  the effective equipment impairment factor,  $H(x)$  the unit step function, and  $A$  the correcting factor, which takes into account the environment where the communication takes place. Besides, ITU-T G.113 (2007) provides a formula to translate the  $R$  value into MOS:

$$MOS = 1 + 0.035 \cdot R + R \cdot (R - 60) \cdot (100 - R) \cdot 7 \cdot 10^{-6} \quad (3)$$

The impairment factors, in turn, depend on the specific codec used for the VoIP communication; the values of these factors for a number of codecs are tabulated in ITU-T G.113 (2007) and its amendment 1 (2009).

### 3 PERFORMANCE RESULTS

In this section, a set of performance figures are shown for web browsing, video YouTube and Skype-based VoIP over a LTE cellular network. Radio performance indicators (at PHY/MAC layers) have been obtained from a dynamic link level LTE simulator (Gómez et al., 2010a), whose main configuration parameters are listed in Table 2.

Average throughput results (per user) at the MAC layer as a function of the received average SNR are shown in Figure 3. Assuming an average SNR of 20 dB, a user would be able to achieve around 4 Mbps considering that the radio resources are shared among 10 users. Regarding the BLER and delay at the MAC layer (not shown in the figure), they have been also obtained from simulations, whose values are:  $BLER \approx 5\%$ ,  $delay_{MAC} \approx 15$  ms.

The following sections will use MAC layer results as a baseline for upper layer performance estimation.

Table 2: PHY/MAC configuration parameters.

Parameter	Value
Carrier frequency	2 GHz
System bandwidth	20 MHz
Duplexing scheme	FDD
Resource block (RB) BW	180 KHz
Subcarriers per RB	12
Sub-frame duration	1 ms
Antenna configuration	1-layer MIMO 2x2
Precoding	LTE 4-words codebook
Power delay profile	Extended pedestrian A channel
UE speed	4 km/h
Channel estimation	Zero-Forcing
MIMO Detection	MMSE
Target BLER	10%
Control channel overhead	From 1 to 3 OFDM symbols
Modulation / coding rate	16 CQI table (4bits)
Coding scheme	Turbo codes + SOVA
CQI & PMI delay	1 ms
CQI reporting period	1 ms
HARQ model	Incremental Redundancy + Chase Combining
# stop and wait processes	8
Scheduling method	Proportional Fair
Averaging window size	500 ms
Number of users	10
Source model	Full buffer

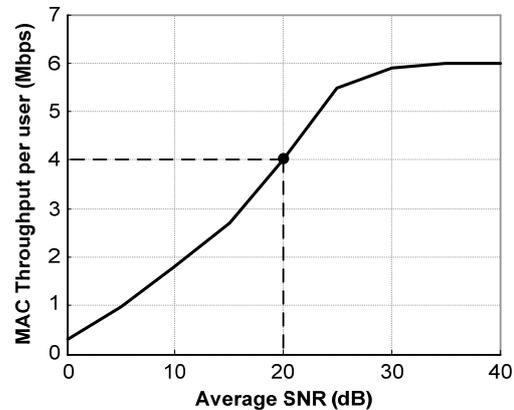


Figure 3: Average throughput at the MAC layer vs SNR.

#### 3.1 Web Browsing

The Web service architecture uses a client-server approach in which the exchange of information is done via HTTP/TCP. HTTP version 1.1 has been assumed in the analysis. This version includes the persistent connection feature, which makes it possible to reuse the same TCP connection for downloading subsequent objects. The pipelining

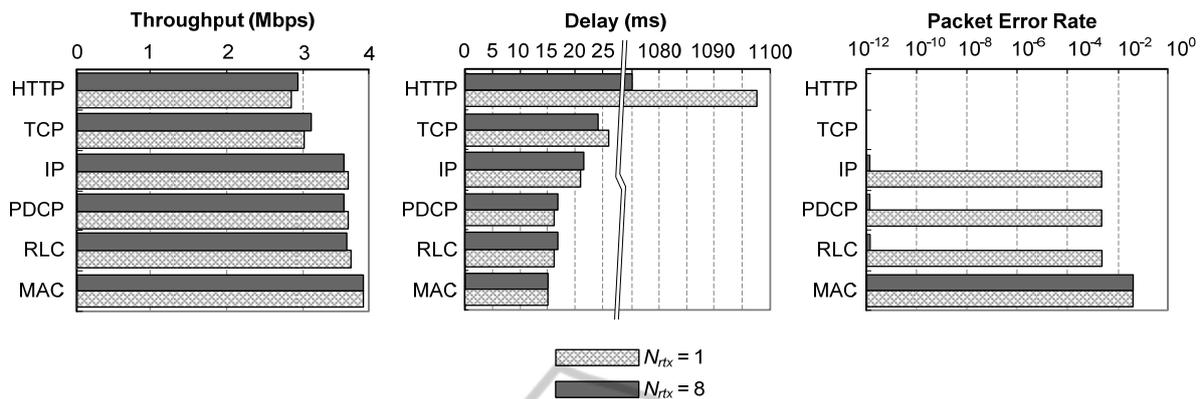


Figure 4: Performance evaluation at each protocol layer (web browsing).

feature have been also assumed, thus allowing a number of object requests to be sent without waiting for the reception of the previous object.

Figure 4 shows a performance analysis at each protocol layer of the web browsing service, starting from the MAC layer results described in previous section. A web page consisting of 100kB text and 15 secondary objects whose individual average size is 20kB has been considered (400kB page size). It is assumed that all the objects and text are located in the same web server, so that all the data transfers will run on top of the same TCP connection.

Regarding TCP configuration, the following settings have been used: maximum segment size  $MSS = 1460$  bytes, initial congestion window  $W_{init} = 1$  segment, advertised window from the receiver  $AWND = 32$  kbytes, number of ACKs per transmitted segment  $b = 1$ , and SYN timeout  $T_s = 3s$ .

Results shown in Figure 4 provide a detailed analysis of the performance achieved at each protocol layer (in terms of throughput, delay and loss rate). Let us analyze the performance in a bottom-up approach, starting from the MAC layer (obtained from simulations) up to the application:

- MAC error rate can be lowered by means of RLC level retransmissions (ARQ protocol). The graph shows the results for two different values of the maximum number of RLC retransmissions ( $N_{rx}$ ): 1 and 8. Results show that higher  $N_{rx}$  values make it possible to reduce the error rate at the expense of decreasing the effective throughput and increasing the delay at RLC layer. However, when TCP is used at transport layer, it is highly recommended to decrease the error rate at lower layers so that end-to-end retransmissions are avoided.
- PDCP layer does not apply header compression in this scenario, so its impact on the performance indicators only comes from the PDCP header overhead.

- At the IP layer, the delay from the base station to the web server is assumed to be 5ms whereas the packet loss rate is negligible compared to the radio interface (the main focus of the analysis is given to the impact of the radio interface on upper layers).

- TCP behavior is very sensitive to IP loss rate, as its congestion control protocol tries to adapt the instantaneous transmission rate to the network characteristics in order to provide reliability, i.e. loss rate zero. In that sense, if IP loss rate is minimized at the radio interface by means of a higher number of local retransmissions (e.g.  $N_{rx} = 8$ ), TCP will be able to achieve higher average sending rates; additionally, in that situation, average TCP delay is also reduced as the number of end-to-end TCP retransmissions is decreased.

- At application layer, HTTP delay results represent the complete “click-to-download” time of the whole web page, including: DNS query, TCP connection establishment, text and secondary objects request and download.

From previous analysis, it is important to highlight the impact of the loss rate on TCP performance. For that reason, the reliability of lower layers is an issue when the radio conditions are poor.

The results associated to the web page downloading time ( $D_{L5}$ ) and MOS, computed from Eq. (1), for different MAC error rate values are shown in Figures 5 and 6. Three different RLC configurations have been evaluated: Unacknowledged Mode (UM), which does not perform any retransmissions, and Acknowledged Mode (AM) with 1 and 8 as maximum number of retransmissions ( $N_{rx}$ ). As shown, the difference between RLC transmission modes increases for higher error rates, being AM with  $N_{rx} = 8$  the best performing configuration since it provides a MOS > 4 (Good) up to 20% of MAC error rate .

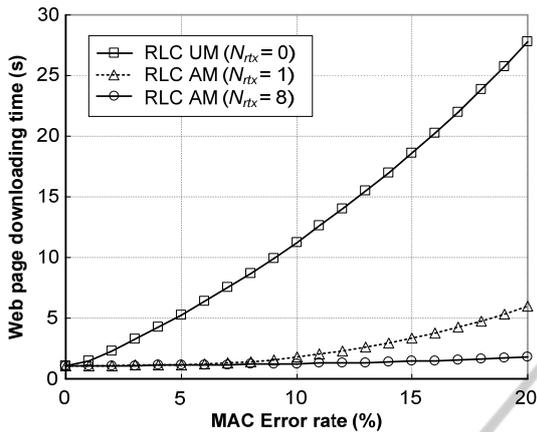


Figure 5: Web page downloading time and MOS.

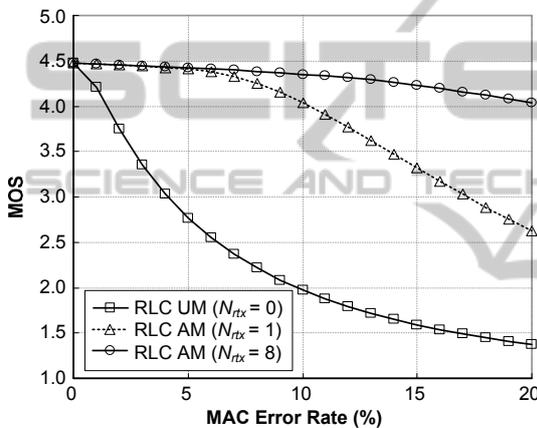


Figure 6: MOS results for web browsing service.

Figure 7 shows the MOS results for different network RTTs and different number of secondary objects in the web page. Firstly, long RTTs lead to a worse TCP performance (in terms of throughput) as a consequence of its inherent congestion control mechanisms based on a transmission window (both during slow start and steady state phases). Such throughput reduction has a direct impact on the web page downloading time and MOS. Secondly, a higher number of objects in the web page (assuming equal sizes) leads to longer downloading times. This behavior may be enhanced by using pipelining feature, which provides higher gains as the number of objects is increased. Pipelining can be achieved to different extents depending on how the request-sending is scheduled on the client's browser. In the figure, a totally pipelined scenario is assumed, i.e. all the object requests are sent in parallel. If a lower number of parallel requests is configured at the browser, the results would be located between both curves (shaded area in the graph). If we compare the results between 5 and 30 secondary objects (i.e.

200kB and 700kB including the text), it can be concluded that much shorter RTTs are required to keep the same MOS (e.g. 110ms and 40ms, respectively to achieve good performance: MOS=4).

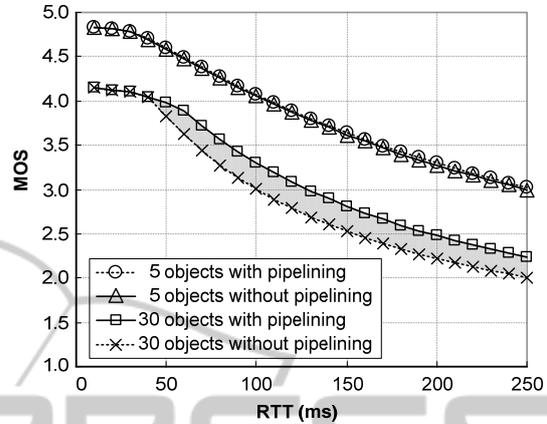


Figure 7: MOS (web browsing) for different RTT.

### 3.2 Youtube

YouTube service is based on progressive download technique, as explained by Gill et al. (2007), which enables the video playback before the content download is completely finished. As data is being downloaded, the video content is temporarily stored in a buffer at the client side, thus enabling the video playback before having the complete video file. This technique is based on HTTP/TCP, i.e. the client sends an HTTP request and, as a consequence, the YouTube multimedia server delivers the requested video through an HTTP response over TCP. The process of downloading the video content from YouTube multimedia server consists of two phases: initial burst (in which data are sent as fast as possible using the whole available bandwidth) and throttling algorithm (in which data are sent at a constant rate related with the video coding rate). Once the video playback has started (which implies that the buffer has certain data to be consumed), if a network congestion episode takes place, the data that are not able to be delivered (from the server) at this constant rate will be later transmitted at the maximum available bandwidth as soon as the congestion is alleviated. This circumstance could trigger a rebuffering event if the client buffer runs out of data. In this case the video playback will be paused until the data buffer is restored. Otherwise, the rebuffering event will be avoided and the congestion will be seamlessly elapsed to the user.

Figure 8 depicts the results of the application performance metrics for YouTube (defined in

section 2.2) as a function of the network RTT for a particular RLC transmission mode (AM with  $N_{rx} = 8$  retransmissions). The following application settings have been used: video length = 250s, client data buffer necessary to start the playback  $B_{full} = 32s$ , buffer threshold that triggers a rebuffering event  $B_{empty} = 2s$ , and video coding rate = 512kbps. Regarding TCP settings, the same configuration as defined for web browsing have been considered.

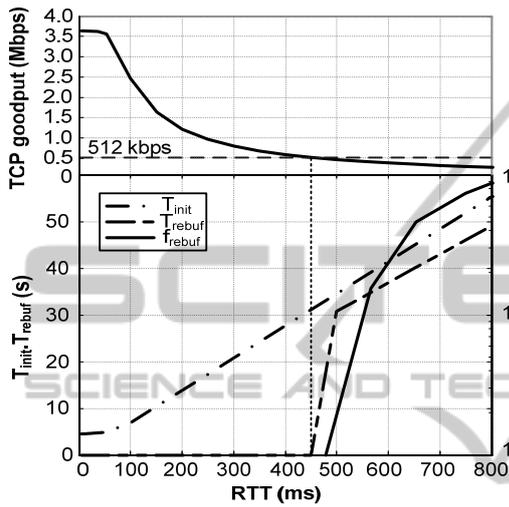


Figure 8: Application Performance Metrics for YouTube as a function of RTT (RLC AM,  $N_{rx}=8$ ).

The upper graph in Figure 8 represents the achievable average TCP goodput (Padhye et al., 1998) for the specified TCP configuration, network RTT and loss rate ( $\approx 2 \cdot 10^{-12}$  as shown in Figure for web browsing). So if the average TCP goodput is higher than the video coding rate (512kbps), then no rebuffering events will take place. As the RTT is increased, the TCP goodput is decreased until it becomes lower than the video coding rate at certain RTT value; from this RTT value and above, the parameters related to the rebuffering events ( $T_{rebuf}$  and  $f_{rebuf}$ ) are higher than zero (as shown in the lower graph). The initial buffering time ( $T_{init}$ ) is also increased for higher RTTs since lower TCP goodput values lead to longer delays to reach the minimum buffer occupancy ( $B_{full}$ ). The rebuffering time ( $T_{rebuf}$ ) has the same behavior, although it is null as long as TCP goodput is above the video coding rate (i.e. no rebufferings occur). Besides, it can be seen that  $T_{rebuf} < T_{init}$  for the same RTT value due to the following reasons: 1) the amount of data needed to be filled ( $B_{full}$ ) for the computation of  $T_{init}$  is greater than the amount of data ( $B_{full} - B_{empty}$ ) required for the computation  $T_{rebuf}$ ; and 2) the computation of  $T_{init}$  assumes that TCP data transfer start with a slow start

phase whereas the computation of  $T_{rebuf}$  considers the TCP steady state to be reached (being the TCP goodput higher in this second phase).

Figure 9 shows the MOS results, from Eq. (2), for different RTTs and RLC transmission modes. As mentioned above, for low RTT values (which achieve TCP goodput values higher than the video coding rate), the initial buffering time is the only metric affecting the MOS (the higher the  $T_{init}$ , the lower the MOS). When the rebuffering events start to take effect over the MOS, its value is rapidly decreased, since interruptions over the playback are annoying for the users. As shown in Figure 8, MOS results could be improved by selecting a proper RLC transmission mode: MOS values are higher for RLC AM mode than for UM mode. It can also be seen that the minimum RTT value that triggers rebuffering events is higher when the AM mode is selected, and even further for a larger number of RLC retransmissions.

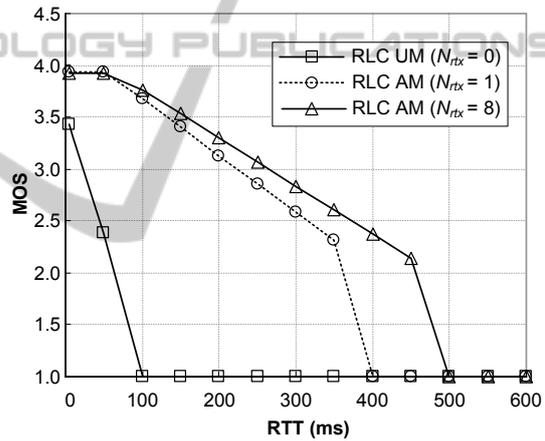


Figure 9: YouTube MOS results for different RTTs and RLC transmission modes.

### 3.3 Skype

In this section, the performance of a Voice over IP (VoIP) service using Skype is analyzed. Skype usually relies on UDP as transport layer, unless the UDP communication is unfeasible, in which case Skype would fall back to TCP. We will focus on the usual Skype behavior over UDP. This section is focused on the E2E communication (i.e. between two Skype clients). The codec used by the software has a big impact on the service performance, being SILK the codec currently supported for E2E communications (since version 4.0). This codec has a set of coding rates from 6kbps to 40kbps. Due to the low data rates that a VoIP flow usually needs, throughput requirements at the network side are not

usually an issue for Skype service. Instead, the network performance indicators mostly affecting the service quality are: loss rate and end-to-end delay.

The following MOS results have been obtained, from Eq (3), for medium and low voice coding rates, considering an A factor value according to a cellular communication inside a building (see section 2.2). The impairment factor associated to this scenario has been obtained from ITU-T G.113 (2007) for the selected voice codecs. Taking into account the characteristics of the VoIP traffic, a Robust Header Compression (RoHC) mechanism has been applied at the PDCP layer. In addition, the RLC UM mode has been selected in order to minimize the end-to-end delay, which is the application layer metric that mostly affects the MOS.

Figure 10 shows the MOS results (for different voice coding rates) as a function of the MAC error rate at the radio interface. As the RLC UM mode has been assumed in this case, potential data errors have a very negative impact on the voice quality. Concretely, fair quality (MOS > 3) is achieved for MAC error rate below 2.5% (for 23.85 kbps) and 1% (for 8.85 kbps). In order to solve this problem, stricter target BLER values are recommended to be configured at the physical layer so that more robust coding schemes are applied.

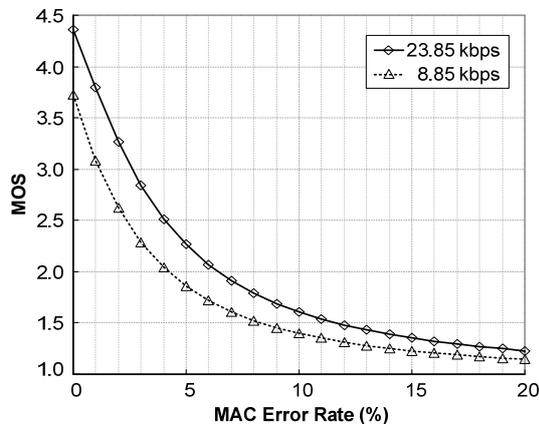


Figure 10: MOS for Skype for different MAC error rates.

Figure 11 shows the MOS results as a function of the end-to-end UDP delay, which has a lower impact on the MOS (for the range of typical delay values) than the error rate. Furthermore, it can be observed that when using default MAC error rate results (5%), MOS results are always poor (i.e. below 3) even for negligible end-to-end delays. If 1% error rate is considered at the MAC layer, maximum end-to-end delays that makes it possible to obtain a fair quality are around 100 ms (for 8.85kbps) and 270 ms (for

23.85kbps).

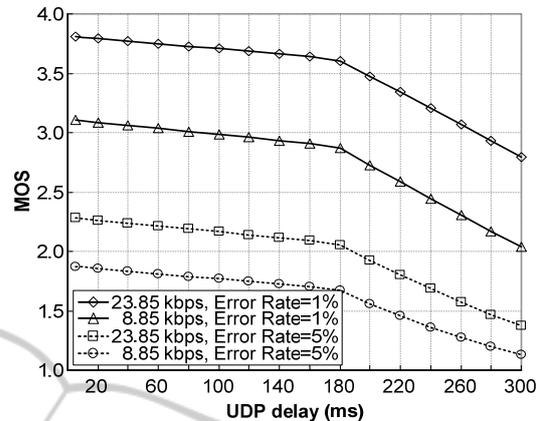


Figure 11: MOS for Skype depending on the UDP delay.

## 4 CONCLUSIONS

This paper presents a QoS and QoE performance evaluation method for data services over cellular networks. In particular, a bottom-up performance analysis have been proposed for evaluating the application layer metrics whereas a set of service-specific utility functions have been used to estimate the MOS for web browsing, video YouTube and VoIP-based on Skype. The methodology here proposed makes it possible to identify the sources of performance degradation along different elements and protocols in addition to the end-users' experienced quality. Additionally, this approach provides the following advantages: 1) it makes it possible to predict the QoE when measurements in real network are not available; 2) it is applicable to any service and wireless network, simply by providing appropriate models; 3) services and networks under analysis do not necessarily require being up and running.

Performance results for web browsing show the great impact of the network loss rate on TCP performance, thus a proper configuration at the radio protocols is a key issue to improve the QoE; additionally, the network RTT is also a critical performance indicator, which subtracts  $\approx 1$  point from the MOS scale with each additional 100ms. In the case of YouTube, results are very dependent on the video coding rate and network metrics (RTT and loss rate); our performance estimations show that fair quality (MOS > 3) can be obtained for RTTs below 200ms when an RLC AM is configured. Finally, Skype results show the great influence of the voice coding rate and error rate on the MOS as

the RLC UM is usually configured for VoIP. End-to-end delay also plays an important role in Skype performance, whose maximum admissible value depends on the coding rate and error rate in the network (as an example, a maximum delay of 100ms is admissible for 8.85kbps and 1% MAC error rate).

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