

# Path Planning Optimization based on Bézier Curves through Open-doors Way Point

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Abstract: Generalized Voronoï Diagrams has been demonstrated to be a relevant tool for planification in a mobile robotics context. Therefore, the generated trajectories may suffer of discontinuities and non-optimality. This paper introduces a reflexion on the use of Bézier curves to solve both of these drawbacks. The key idea of this paper is to be able to smoothen a trajectory in order to save traveling time and therefore reduce displacement and overall consumption (in our mobile robotics context: reduction of battery usage and localization errors). The presented work is firstly detailed and explained on a synthetic map, and experimental results with mobile robots are presented. Disadvantages and advantages are discussed at the end of the paper.

SCIENCE AND TECHNOLOGY PUBLICATIONS

## 1 INTRODUCTION

Path planning is a key task in many fields, especially in mobile robotics. It started in the early 60' with the first industrial robot. Nowadays, we can find those applications everywhere, from the industrial application with the robotic arms assembling cars, to the personal house cleaning "Roomba" robot. However, the solutions are different, the movement is not developed in the same way. The first one is developed under the supervision of a human operator, to guarantee a correct accuracy and be sure of the repeatability of the placement in a defined universe. In the second case, the trajectory is studied "on-line" and completely autonomously. The robot will move in the space and design a map of the environment to determine the path to follow (Jagannathan et al., 1994) and (Dierks and Jagannathan, 2009).

In our context, we had to develop a robot able to discover and create a map of it on-line. Therefore, it needed to plan the trajectory and move around efficiently simultaneously. It is in these conditions that this solution as been submitted as a research work.

The aim of the present work is to improve the current planification and to be able to have an efficient and smooth trajectory. Those two characteristics are important for different reasons. The first is, of course, the traveling time. The more direct is the trajectory and the more time is saved. This result has consequences on other levels. For example, the less the robot travels and the less energy will be consumed

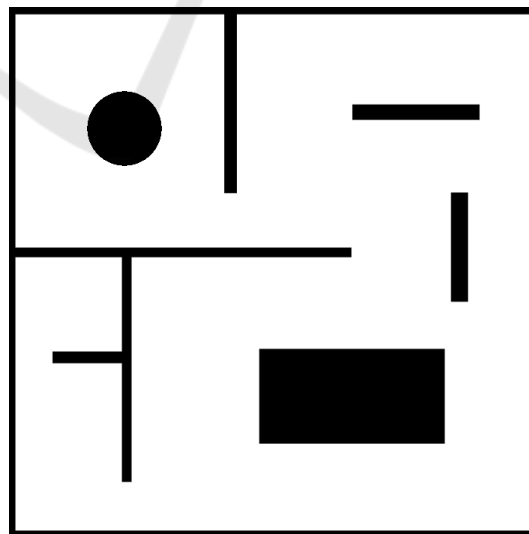


Figure 1: Map used for the demonstration.

and the fastest it travels the more time remains to increase the exploration.

This paper will now introduce the different optimization and algorithms used to develop the trajectory. Some algorithms are already known. They will be explained here as a state of the art and why they are useful in our application.

Note that the demonstration will be supported by examples applied on a synthetic map. In order to have comparable results, the same map will be kept all along and can be seen on the Figure 1.

Let us assume the map is discretized and stored in

the computer's memory as a grid of cells. Each cell represent an information about the space (typ. obstacle or not). For this study, we will store a Boolean in each cell meaning the presence of an obstacle (true) or not (false).

Working straight from the discretized map generates the first problem : the space size for the research of the best path. In two dimensions, we will have a  $n \times m$  cells. If both dimensions grow by 2, the size of the research space is then multiplied by four. The growth is exponential with the dimension.

Another problem with this representation is the accuracy. A fine representation of the environment and path needs a high sample rate that increases the size of the overall grid, memory consumption and computing time.

A last issue with this representation is the graph of the connection between each node (generally equivalent to cells). The bigger the space and the bigger will be the number of connection to pile up in memory. For a 2 dimensional space, there are different kinds of connectivity to define links between cells: 4-connectivity (only horizontal and vertical moves) and 8 connectivity (same as before, moreover diagonal moves are allowed). In a mobile robotics context, none of them, combined with classical approaches (Dijkstra or A\* based algorithms) provides a satisfying trajectory.

Some other techniques like rrt (LaValle, 2006)(rapidly random explored trees) has also been proposed in the state of the art. Even if these techniques has been proven to be very fast and quite efficient in practice, some drawbacks still remain. Such algorithms provides non-optimal solutions without any guaranty of convergence. As these techniques are based on a random exploration of the environment, the repeatability is very poor. For these reasons, our work is mainly focused on a Voronoi Diagram based approach that seems to be an interesting avenue for research.

Previous works have shown that the use of Generalized Voronoi Diagrams (Fortune, 1987) (Seda and Pich, 2008) is a big step forward to tackle those problems. Such diagrams are used to quickly explore the space and thus reduce it to a graph where nodes are particular cells. Indeed, each cell of the graph represents a point which is equidistant to any change in the space (in our case, a change will be represented by the configuration of each cell: navigable or not).

Applying this representation to a standard map will result in the obtention of a graph with nodes, connected to each other. Nodes are only placed in the middle of the free space between walls. Such computed diagram is shown on Figure 2.

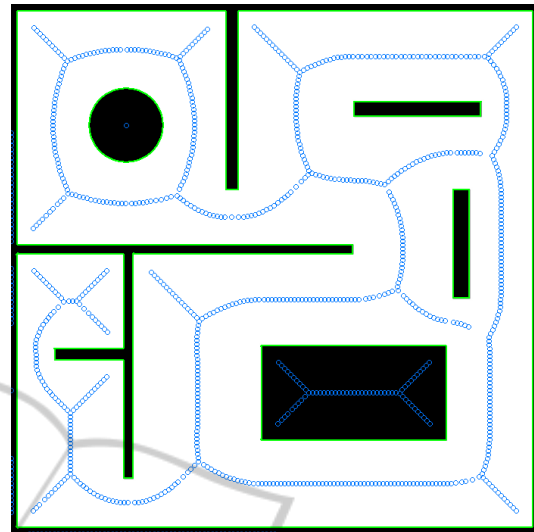


Figure 2: Voronoi graph on a map (the blue circles are the Voronoi cells).

There is still one limit which remains similar to the one before: the granularity of the space. The finer the graph will be and bigger memory space will be used. However, this amount will always be very smaller than the initial grid representation. Note also that the computation of the Voronoi Diagram has been proven to have a linear complexity (Fortune, 1987).

## 2 TRAJECTORY PLANNING

To simplify our presentation, we will consider that the graph is connected, from one point we can always reach another one. The input of our path optimization algorithm can be obtained by any path planning algorithms (A\*, Dijkstra algorithm, ...). On Figure 3, the path from the top left node of the Voronoi graph to the bottom right was computed using the Dijkstra's algorithm.

### 2.1 Optimization Issue

As seen on the Figure 3, different spots can be optimized. For example, the orange circled area is not optimal. Indeed, to be more efficient, the trajectory would have to go less high and more straight. In this way the robot would save energy (less turn and acceleration) and use less time to travel.

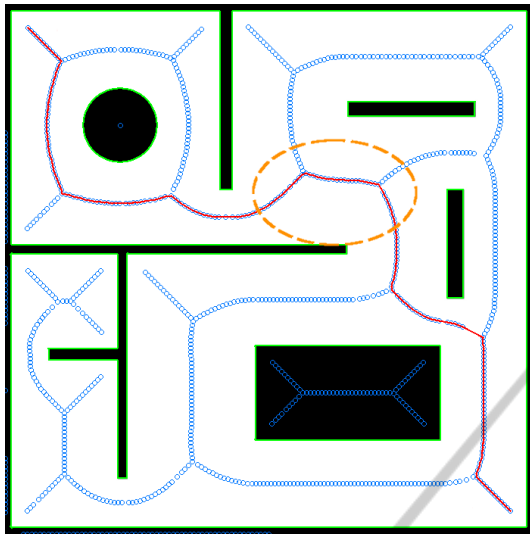


Figure 3: Voronoi trajectory with a possible optimization area.

### 3 BÉZIER CURVE BASED TRAJECTORY OPTIMIZATION

#### 3.1 Bézier Curve

In order to solve the problem described in the previous section, the trajectory need to be smoothed. To do so, the application of a polynomial Bézier (Demengel and Pouget, 1998) curve can be used. In a simple way, this curve will use each node from the Voronoi graph which are visited by the path finding algorithm to be computed. A node will then be considered as a "way-point" of the Bézier curve. To draw this curve, each way-point is pondered by a polynomial coefficient at a time  $t$ . The definition function of this curve will be:  $\sum_{i=0}^n B_i^n(t) \cdot P_i$  with  $t \in [0, 1]$ ,  $B_i^n$  are the Bernstein coefficient and  $P_i$  are the way-points.

The problem with this approach is that the number of way-point will decrease the smoothing effect of the curve. In order to counter that, another approach is followed, using less points and a "part-to-part" definition of the curve.

#### 3.2 Gate Way-point

To avoid the lost of the smoothing effect of the Bézier curve by using too many way-point, a new method of way-point definition is here proposed. This algorithm is based on the human behavior when going trough a door. The main idea is that whenever we (humans) want to go through a door, we will do it with our shoulder oriented in the perpendicular axis of the door, and the body centered in the doorway.

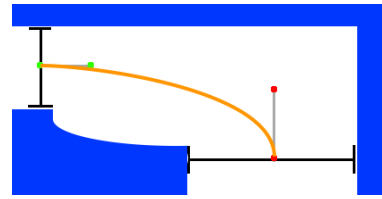


Figure 4: Bézier curve (in orange) through doors (green and red dots are Bézier way-point for each door and blue area are walls).

Another observation point is that in a space, a door represent a local minimum in the function defining the trajectory. Therefore, it is rather easy to detect them.

From this statement, it is decided that the center of all the "doors" will be a way-point for the Bézier curve. To be able to compute a Bézier curve, at least 4 points are needed. From one door to another it makes only two. Two others will be added upon the direction and the size of the door. Two way point are defined in the following way:

- the way-point is placed on a virtual line perpendicular to the axis of the door,
- the wider is the door and the furthest from its center can be the way-point.

From this stage, four way-point are defined. The Bézier curve can be defined using the following expression:

$$\mathbf{P}(t) = \mathbf{P}_0(1-t)^3 + 3\mathbf{P}_1t(1-t)^2 + 3\mathbf{P}_2t^2(1-t) + \mathbf{P}_3t^3$$

for  $0 \leq t \leq 1$  as seen on Figure 4.

To define a complete path, we just need to define a trajectory from door to door and then merge all those path together to obtain the complete navigation. A big advantage of this method is that the navigation will always be safe when going through critical places (the doors, local minimum), as seen on Figure 5. However, the method is not yet proved to be safe (obstacle-free) in between. The next part will demonstrate why and what solutions/tests could be experimented.

### 3.3 Drawback, Limitation and Discussion

#### 3.3.1 Size of the Door

As seen on the Figure 5, one of the first limit of this Bézier smoothing is that the doors are not limited in width. It means that even in a big space there is going to be a door to start it and so a way-point for the curve. The result is that the curve as to make some rather big detour and so make a lost of time for the robot.

One way to sort out this limit is to use a threshold on the size of a door. If it's bigger than this threshold, it is not taken in account to create a way-point.

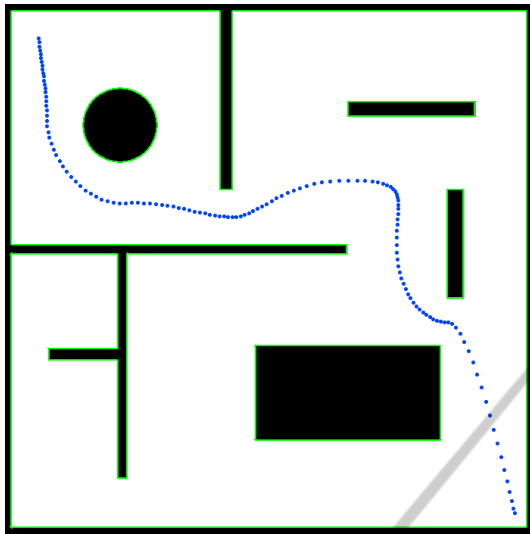


Figure 5: Bézier curve (in blue) through doors.

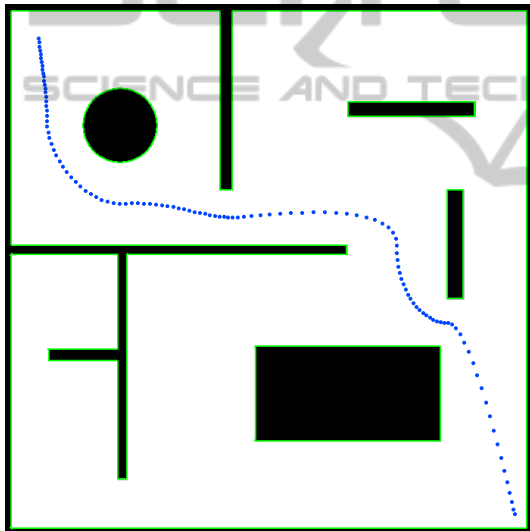


Figure 6: Bézier curve (in blue) through doors, with a limit on the "door width".

However, this solution depends on the context (size of the corridor, scale of the map...) and is difficult to be determined in an autonomous way of decision. A manual thresholding could give the result seen on the Figure 6 as opposite to the Figure 5 without it.

### 3.3.2 Safety of Traveling

As for now, the question of safety of traveling cannot be answered (as stated in the introduction, it is still a work in progress). The study is now at this point of solving the answer of the question: "Is it always collision-free?". In most of the case, the empiric results show that the computed smooth trajectory is safe. However, in some particular case (with aligned

circular wall) the path will go *through* the wall (See figure 7).

As a perspective of this work, two solutions have been emitted. First, a property of the Bézier curve says that the curve is always contained in the convex hull that defined it (See figure 8). This mean that we could check if the convex hull intersects a wall or not. If yes, then the trajectory between those two doors need to be improved. If no, then this part of the curve is collision-free.

Another idea is to check if the circumscribing circle of the envelope is within a corridor defined by the Voronoï graph. The notion of the corridor is given by the union of all circular spaces between a Voronoï site and the obstacles it refers to (as seen on the Figure 9). In this circle, no obstacles can be found. Therefore, testing if the circumscribing circle of the envelope is in the corridor allows to conclude on the safeness of the path. The advantage of this solution is the speed of the computation; it is quite fast to check if a circle is included in an union of other circles.

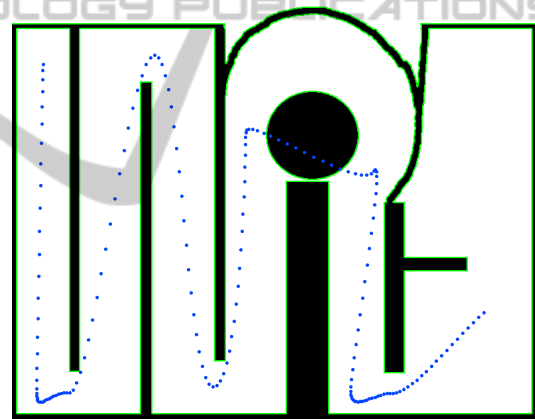


Figure 7: Non-obstacle-free trajectory.

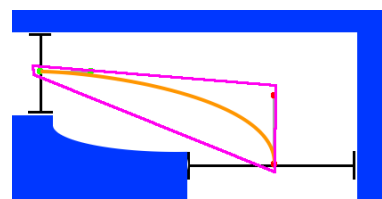


Figure 8: Envelope (pink) of a Bézier curve (orange).

## 4 EXPERIMENTS

This work takes place in a larger project called Cart-O-matic. Our team was involved in a robotics competition (Défi-CAROTTE) founded by the French Research Agency (ANR) and the General Delegation for Armaments (DGA). The aim of this contest was

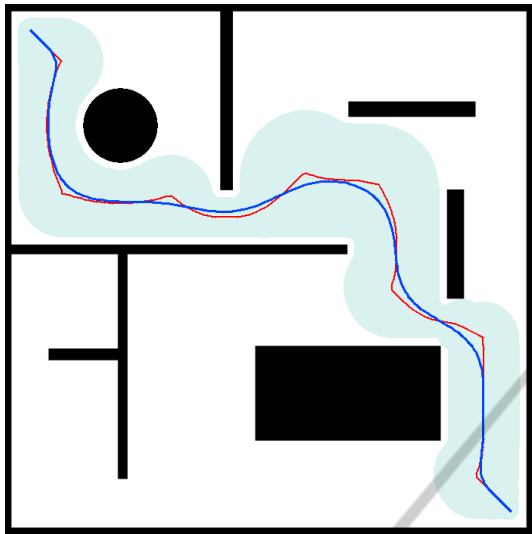


Figure 9: The corridor (Grey) for a path (blue).

to map and locate objects in a structured environment similar to an apartment. The particularity of our team was the use of a multi-robot strategy (Shahbandi and Lucidarme, 2012) (Bautin et al., 2011). Our team designed and built seven identical mobile robots called MiniRex (MINIature Robot for Exploration) illustrated in Fig. 10. Each robot is composed of an Embedded PC (proc. Atom 1.6GHz), inclinometer, ultrasonic sensors for navigation, LIDAR for localization and mapping, and an RGB-D sensor (Microsoft Kinect) for object recognition. Figures 11 and 12 illustrate the proposed algorithm applied on environments mapped by robots.



Figure 10: The MiniRex robot while exploring its environment.

## 5 CONCLUSIONS

As seen in the development, this optimization method introduces advantages on the path planification problems. The global idea of reducing the best trajectory seems to be reached and the "human-based" behavior

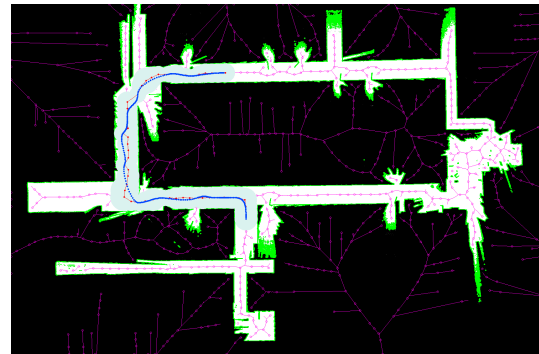


Figure 11: Illustration of a trajectory inside a mapped building from the university of Angers.

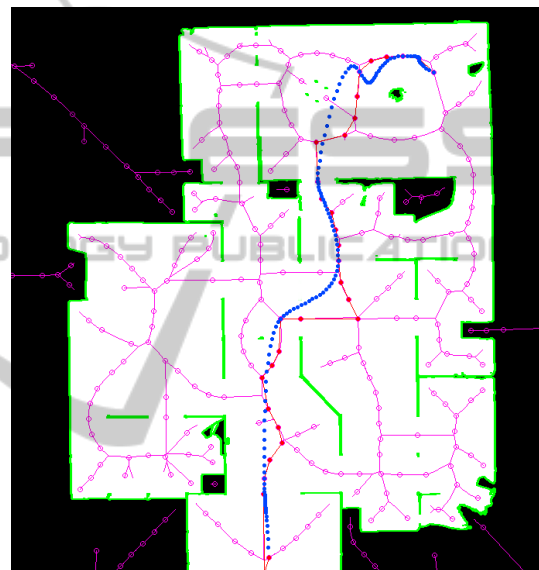


Figure 12: Illustration of a trajectory in a multi-robot mapped building.

tends to give a reliable solution and elegant way of displacement.

However, even if in most cases the algorithm seems to work, the lack of a mathematical proof cannot allow to conclude on the efficiency of the method. Moreover, as seen on Figure 7, some cases brings a set of new problem to which solutions are not found yet.

## ACKNOWLEDGMENTS

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