A Simulation Framework for Analyzing Complex Infinitely Repeated Games

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We discuss a technique to analyze complex infinitely repeated games using techniques from the fields of game theory and simulations. Our research is motivated by the analysis of electronic markets with thousands of participants and possibly complex strategic behavior. We consider an example of a global market of composed IT services to demonstrate the use of our simulation technique. We present our current work in this area and we want to discuss further approaches for the future.

1 INTRODUCTION

The field of game theory has developed concepts and methods that describe and predict the outcome of strategic situations in which rationally acting individuals interact. However, when analyzing large-scale complex systems these methods quickly reach their limits. Even for games without repetition and even in games with few players the problem of determining the outcome is computationally difficult (Daskalakis et al., 2009; Chen et al., 2009). In games that are being played repeatedly the situation is even more unsatisfactory as the behavior of players is a lot more complex than the mere choice of an action. This makes the analysis of such games a daunting task even if one imposes severe restrictions on allowed strategies for the players. Therefore, we choose a simulationbased approach to analyze settings of repeated strategic interaction in complex scenarios. To that end, we model the strategic behavior of players as algorithms or automata which allow us to simulate the interaction of many players with possibly many different strategies. The results of the simulations provide us with information about the performances of strategies and which of them are plausibly chosen in real-world scenarios. Our approach nicely complements experimental as well as theoretical methods.

Our research is motivated by the analysis of electronic markets with thousands of participants and complex strategic behavior. We consider an example of a global market of composed IT services (Happe et al., 2013) to demonstrate the use of our simulation technique. In this setting service providers offer simple software services on an electronic market platform. These services can dynamically and flexibly be combined to more complex and individual service compositions by service composers. An important characteristic of this market model is a certain degree of anonymity and information asymmetry. Service providers may have an incentive to exert low effort in providing their service, resulting in lower production cost and lower quality. The service composer would have to predict a service provider's choice of effort although he might not even be able to observe it in hindsight. We seek to understand the impact of this asymmetric information on the quality of the services and how it can be improved by market mechanisms like reputation systems (Brangewitz et al., 2014).

We explore the capabilities and limitations of such simulations by performing an exhaustive search over the strategy space of repeated games with finite automata, also known as machine games (Rubinstein, 1998; Rubinstein, 1986). Similar evaluations have already been investigated (Axelrod, 1984; Izquierdo and Izquierdo, 2006; Pence and Buchak, 2012), but they all concentrated on simple game scenarios like the Repeated Prisoner's Dilemma. In the following we present our underlying model and the simulation framework before we discuss the results of several simulated scenarios.

Abstract:

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2 MODEL

We consider the model of an infinitely repeated game $\mathcal{G} = (N, Q, (A_i)_{i \in N}, \delta, (u_i)_{i \in N})$ consisting of a set of *n* players $N = \{1, ..., n\}$, a set of states Q, for each player $i \in N$ a set of actions A_i , a transition probability function $\delta: Q \times A_1 \times \ldots \times A_n \times Q \rightarrow [0,1]$ and for each player *i* a utility function $u_i: Q \times A_1 \times \ldots \times A_n \rightarrow Q$ \mathbb{R} (Luce and Raiffa, 1957; Aumann, 1959; Shapley, 1953). The strategic behavior or strategy of a player is a function or algorithm that maps the history of the game to the player's next action. Which information about the history each player observes, how much memory he possesses and which computations he is able to perform is left open here. Thus, this definition captures simpler models like positional games or stochastic games as well as more complex models like machine games (Rubinstein, 1998; Rubinstein, 1986).

We model the introduced market scenario as such a repeated game that proceeds in rounds. In each round a service composer is asked to compose a certain service product. That is, he has to choose a set of services (out of a collection of possible choices) and composes the final product. The service providers which have been chosen by the composer can decide how much effort they exert in delivering their services. For simplicity we assume only two possible levels, high and low. The utility of a service provider is 20 if he only exerts low effort, 10 for high effort and 0 if he is not chosen by the composer. Each service composition consists of multiple services and, thus, its quality depends on the aggregated service qualities. Here we assume that the quality of a service composition is good if and only if each service was delivered with high quality. The realized quality of the composition can be correctly observed with a certain probability, which we chose to be 95% in our simulations. We implicitly assume that a player's reputation depends on the observed quality of the services he participated in during the past rounds. We assign to each player a reputation vector of fixed length (in our example it will be two or three) storing the past (two or three) observed qualities. Note that the reputation depends on the choices of all participating service providers and the outcome of the random event of the observation. The payoff of the composer depends on his reputation and is 40 plus (minus) 10 for each positive (negative) entry in his reputation vector. The overall payoff of the repeated game is the discounted reward with discount factor of 10% for each player.

In general, the strategy of a player is modeled by a deterministic transition function $s_i : Q \to A_i$ which determines the player's next action, depending on the current observed state. Here, we assume that a player's observable state is only his last reputation values. Thus, the class of transition functions is restricted to functions that do not differentiate between states with the same reputation values.

3 SIMULATOR

3.1 Simulation Architecture

The simulator is a packaged Java application. This allows an execution on nearly every available execution environment nowadays without concerning any system-specific properties. In essence, the application consists of two modules (cf. Fig. 1).



The first one is responsible for the execution of the simulation. It consumes the input and starts the simulations with the correct settings in the *Runner* and the *Configurator* components. As input an XML description of the scenario together with further individual properties yields. Furthermore, a *Simulation* and a *Round* component handle all simulation- and round-relevant aspects. A *Listener* component is responsible for collecting results and produces *csv* output files with relevant data like the players' utilities.

As the second module we have the concrete application scenario, in our case the market for composed services. The most important components here are the *Scenario*, the *Player* and the *Strategy*. This last module can be replaced by any other scenario implementation, for example the well-known and extensively investigated Repeated Prisoner's Dilemma (Axelrod, 1984) or a completely new designed scenario.

3.2 Simulation Workflow

The simulation is entirely integrated in an automated workflow. Beginning with an input all required processing steps are executed until the output is generated (cf. Fig. 2). The input consists of four parts:

- a general *configuration* which defines the scenario and all fixed parameters given as XML file,
- a list of *parameters* from which all combinations are evaluated in comparison,
- a list of possible *strategies* for each player in the scenario from which all strategy combinations are evaluated, and

Players	2 reputation values			3 reputation values			
	Combinations	Memory	Runtime	Combinations	Memory	Runtime	
2	256	2MB	30.17sec	65,536	0.32GB	0,09h	
3	4,096	21MB	232.18sec	16,777,216	96.07GB	8,54h	
4	65,536	375MB	344.19sec	4,294,967,296	24,594.88GB	2186,37h	

Table 1: Runtime and memory complexity of evaluations (The numbers written in italics are predicted).

	Table 2:	Four evaluated	scenarios	with 2.3	3 and 4	determinist	ic strategic p	lavers.
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Scenario	Composition	Service Provider	Players	Reputation
1	composition of all	2 deterministic players	2	3 values
2	choose best one	2 deterministic players	2	3 values
3	choose best pair	3 deterministic players	3	3 values
4	choose with one monopolist	3 deterministic players	4	2 values



Figure 2: The simulation workflow.

• a list of additional *settings* for the simulation, i. e., the number of random seeds, the connection parameters for external clusters, etc.

The preprocessing step reads the inputs and generates all simulation configurations necessary for independent simulations (cf. Sec. 3.1). The execution step produces execution packages, distributes them on different available environments (see Sec. 3.3) and collects the results as soon as they are finished. Afterwards, the postprocessing step consumes the output of the different simulations and aggregates them according to different aspects. In the end, it produces tables with the measured values (*.*csv*), ready-to-use graphics for presentation (*.*png*), and papers (*.*pdf*) and generation scripts (*.*plt*) to manually adapt the graphics with gnuplot (Williams and Kelley, 2012).

3.3 Simulation Environment

The simulation environment consists of two parts: On the one hand, the application running on a dedicated virtual server and, on the other hand, different external machines for running the simulations (cf. Fig. 3).

The application itself has two parts: a front-end managing the user input and output as well as a backend responsible for the simulation workflow. The

Figure 3: The available simulation environment.

front-end is running as web application based on the Java *Spring Framework*. Therefore, the user can handle the simulations through a responsive HTML5 web interface from any device and has not to touch more technical parts. The back-end based on the *Spring Batch* manages the simulation workflow and is especially responsible for the pre- and postprocessing of the simulations. For this purpose four different approaches are implemented:

- VM: The simulations run on the same virtual machine as the application itself in different threads, therefore a scheduler is integrated.
- **Pool:** The simulations run on other desktop machines available in a pool. Also for this purpose, a scheduler exists in the application and the transfer of data is managed via the SSH protocol.
- **HTC:** The simulations run on different virtual machines in a High Throughput Cluster. An existing front-end can be accessed via SSH and is responsible for the scheduling.
- **HPC:** The simulations run on nodes in a High Performance Computing cluster. Also for this purpose a front-end exists to allocate the computation time. While the first two approaches are completely

self-developed and independent of any environment,

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Figure 5: Scenario 1: The payoff of two service providers.

the latter two use existing managing approaches available in computing centers. They offer the advantage that the available computing time and memory is much higher and allow clearly more simulations in the same time.

EVALUATIONS 4

In this section we investigate the capabilities and limitations of our tool by evaluating the market scenario. We start by investigating the possible scenarios for evaluation. Concentrating only on one strategy profile, i.e., one combination of strategies, the simulation framework allows very complex scenarios.

However, for the analysis of strategic games and their outcomes we have to take into account all possible strategy profiles that may be chosen by the players. Traditionally, game theory is interested in equilibria which are strategy profiles in which every player has chosen his optimal strategy given the choices of all other players. As the number of possible

strategies grows exponentially with the size of the au-

tomata and, moreover, the number of strategy profiles grows exponentially with the number of players, evaluating all strategy profiles by simulations quickly becomes intractable (cf. Table 1). Already by using only 4 strategic players and a state space with 3 history values, running on a cluster with 100 available parallel computing nodes can not be simulated with reasonable time and space input. Thus, we first demonstrate our tool with a few small examples before we discuss further research beyond a mere exhaustive search.

Therefore, we consider 4 different scenarios with 2, 3 and 4 strategic players (see Table 2). Scenario 1 and 2 contain two service providers acting according to a deterministic transition function. In the first scenario, the composed service contains both separate service components, whereas in the second scenario only the best service (regarding the past reputation values) is included in the product. In the third scenario, three deterministic service providers are available and the composition consists of the two best services. The last scenario is an extension of scenario 3. There are 3 deterministic service providers, but the





Figure 9: Scenario 3: The payoff of three service providers.

composition is also chosen by a player. Furthermore, one of the providers is fixed (it has a monopoly) and the choice is only between the other two.

All simulations are run for 100 rounds with 20 different initial seeds to get usable statistic results. In the following plots we compare the quality and the discounted utility of different providers together with the possible improvement factor of the first player visualized by the color. The incentive to change the strategy is higher in the brighter areas. In the first scenario in which both services are used for the combination we see two dark areas in the lower left and upper left corner of the plot (cf. Fig. 4). The service provider cannot improve its utility by serving higher quality. Furthermore, we can see in Fig. 5 that the player has all equlibria with high utility values.

In contrast, scenario 2 has another result. Here, always playing with a bad quality is not the best choice for the player (see Fig. 6). He has to cooperate and serve a better quality to be chosen. Furthermore, the equilibria are not with high utility values, states where no improvement can be reached are also found in the lower utility area (see Fig. 7). In the third scenario

Figure 11: Scenario 4: The quality of two service providers.

with three players we see similar results (cf. Fig. 8 and Fig. 9). Multiple areas with equilibria exist where no improvement for a single player is possible. A better analysis is possible by using 3D visualization tools with user interaction.

In the last scenario a monopolist and two other service providers are put together into compositions. We see a difference between the stable strategies of the monopolist (cf. Fig. 10) and the other providers (cf. Fig. 11). The monopolist is always in a nearly stable situation, not depending on its served quality. However, the other provider has only small stable areas.

The simulation framework helps us to study repeated games where the theoretic analysis is too complex and experimental data is not available. We can identify equilibria and stable strategies which are a plausible explanation of players' behavior. For the design of the electronic market and supporting market infrastructure (e.g., reputation systems), we can evaluate possible solutions in various different scenarios.

5 CONCLUSIONS AND FUTURE WORK

Not surprisingly, an exhaustive search over all strategy profiles quickly becomes intractable - even if one considers only strategies implemented by automata with very few states. Nevertheless, simulations give us interesting insights about the outcomes of strategic behavior in complex scenarios like the presented electronic market. By identifying equilibria and strategies with high utilities and low incentive to change, useful statements about the development of markets can be done and the impact of mechanisms like reputation systems can be studied before their implementation. However, we are interested in even more realistic scenarios which allow players with more memory and more complex functions including randomization. A promising direction is the use of techniques from evolutionary game theory which uses a population of strategies, each additionally equipped with a fitness value. Utilizing replication and imitation dynamics (Weibull, 1997; Taylor and Jonker, 1978; Hofbauer and Sigmund, 1998), the population and the fitness values evolve over time. The concept of evolutionary equilibria (Hirshleifer and Rubin, 1982) and evolutionary stable strategies (ESS) (Smith and Price, 1973) describe states and strategies which are (approximately) stable with respect to the aforementioned dynamics. The simulation architecture presented here is easily adaptable towards such concepts and the computational results show that hundreds or even thousands of such experiments can be conducted in a very short time.

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