SIM

A Flexible, Scalable and Expandable Simulation Platform Applying to Lunar Orbit Rendezvous Mission

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Abstract:

In this paper we propose a new simulation platform called SIM, for analyzing parallel and distributed systems. This platform aims to test parallel and distributed architectures and applications. The main characteristics of SIM are flexibility, scalability and expandability. SIM is about five functions: model management, experiment management, distribution management, operation management and node management. To improve the efficiency of project development, new models are designed for lunar orbit rendezvous mission to apply the simulation platform. Finally, a validation process and evaluation tests have been performed to evaluate the SIM platform and lunar orbit rendezvous mission models. The simulation platform and models will lay the foundation for the more validations of autonomy technology in manned lunar landing research.

1 INTRODUCTION

Nowadays as an important research in aerospace field manned lunar landing is getting more and more attention from many space powers such as America, Russia, and other European countries (Bocam et al., 2005 and Santovincenzo, 2004). Apollo is the most representative in the history of exploring the moon. Lunar orbit rendezvous technology has proved to be very useful in the Apollo plan (Reeves, 2005).

The methodology of exploring the moon can be classified into earth orbit rendezvous and lunar orbit rendezvous. The earth orbit rendezvous method separates the spacecraft into many modules which are assembled on the orbit near the earth. As the mainstream of exploring the moon, the earth orbit rendezvous method is widely adopted by many countries such as the Crew Exploration Vehicle of America (Raftery and Fox, 2007), the Flier plan of Russia, and the Architecture Study for Sustainable Lunar Exploration by ESA CDF Study Academy (Santovincenzo, 2004). The earth orbit rendezvous is a mature and highly reliable technology, but it needs more fuel and longer rendezvous period which is very unbeneficial to the whole mission. The lunar orbit rendezvous method is to assemble all the modules of the moon exploring spacecraft on the lunar orbit. Because of the smaller gravity of the moon, it costs less fuels to dock on the lunar orbit. Thence, it is appreciated by more and more countries as a new method to explore the moon.

Lunar orbit rendezvous is an independent and complicated mission. It is necessary to verify the reliability and security of the mission by the technology of distributed simulation. Many countries have committed a large number of resources to build suitable simulation platforms in the related field. These platforms are playing important roles in the ground experiment.

Some of these simulators are focused on simulating. And the entire system provides functional execution of unmodified commercial operating systems and applications such as COTSon (Argollo et al., 2009), a simulator framework jointly developed by HP Labs and AMD that provides accurate evaluations of current and future computing; M5 (Binkert et al., 2006), which supports the execution of the entire system, including operating system code, models of network and disk devices; Simics (Magnusson et al., 2002), another full-system simulator that was one of the first academic projects in this area and the first commercial full-system simulator; and SimOS (Rosenblum et al., 1995 and Rosenblum et al., 1997), an environment for

studying the hardware and software of computer systems. These simulators are called full-system simulators. The main advantage of those simulators is the high level of accuracy obtained, whereas the main drawback is its performance, which in most cases is five or six orders of magnitude slower than a real system.

Moreover, there are approaches that do not focus on modelling and simulating the system with a full level of detail instead on balancing the level of detail to model the system with the performance and accuracy obtained. For instance, Phantom (Zhai et al., 2010) proposes a novel approach to predict the sequential computation time accurately and efficiently by integrating a computation-time acquisition approach with a trace-driven network simulator. dPerf (Cornea and Bourgeois, 2010) is a tool that uses Rose (Liao et al., 2009) for performing static analysis of the input source code of programs written in C, C++, or Fortran.

There are also other works that focused on distributed storage architectures. One example of this kind of system is Modeling Infrastructure for Dynamic Active Storage (MIDAS) (Tarapore et al., 2008). MIDAS is an execution-driven simulator that captures both the processing and I/O behavior of active storage systems. MIDAS simulates a host system interacting with the I/O path via an interconnection network. The simulated I/O path can include disk drives with programmable processors programmable storage controllers. micro-architecture of each one of these components is configurable. With this framework, the effects of different processor micro-architectures, physical disk and network designs, and communication protocols on application performance can be explored.

Due to the high number of domains in the field of distributed systems, developing a universal simulator is impractical and unfeasible. Naturally, each researcher has its own objectives and requirements, and the same way each simulator is developed for a specific purpose. Many existing simulators do not fit the researcher's requirements. As a result, researchers have to modify an existing simulator, or coding a new one. But coding a simulator from scratch is a very complex and difficult task. Usually, researchers use simulation frameworks for building a specific simulator.

In this paper, we propose a new simulation platform called SIM, which is oriented towards analyzing and studying parallel applications on distributed systems. SIM has been designed to provide flexibility, accuracy, performance, and scalability. Those features make it a powerful simulation platform for designing, testing and

analyzing both actual and non-existent architectures. Simulation Systems range from a single computing node to a complete high performance distributed system. In fact, this simulation platform has been applied to data systems simulation in the 921 Manned Space Office of China.

The rest of the paper is structured as follows. Section 2 presents some requirements. Section 3 describes the basic architecture of SIM. Section 4 shows the strategies and the tools to model distributed environments in SIM. Section 5 presents practical implementation and experimental results. Finally, Section 6 presents some conclusions and future works.

2 **REQUIREMENTS**

Actualization of the lunar orbit mission puts forward higher requirements of the project system such as higher precision of the lunch vehicle operational accuracy, more powerful relative navigation or rendezvous and docking of spacecraft, shorter response period of measurement and control communication system, higher precision of measurement and control instrument. Aspects needed to be verified from the whole project are:

- The Mission Profile Verification
 Verifying validity and rationality among the systems and mission phases of the lunar orbit rendezvous.
- (2) Mission Software Verification
 Verifying validity of the software used by the
 lunar orbit rendezvous test experiment. These
 softwares include lunch window calculation,
 orbit determination and fuel injection, spacecraft
 GNC.
- (3) Flight Control Strategy
 Verifying validity of the flight control strategy.
 Verifying the effects of orbit error on the flight control. Verifying the strategy for the orbit fault-pattern. Verifying the optimal methods of the flight control.
- (4) Visual Presentation for Flight Process Visual presentation for the whole flight process of lunar orbit rendezvous. Providing visual image of 3D scene, subastral point of the flight process.

Considering common problems of simulation platforms for different kinds of aerospace missions, we must understand the structure of the new platform and relationship among function layers before designing. We must ensure sufficient versatility, standardization and extendibility of the

platform. The following points shall be paid attention to:

- (1) Most platforms are suited to only one kind of models. In the face of complicated aerospace projects, modeling development is not complicated by one person or one company. Type of singleness brings much trouble to the system integration job.
- (2) Most platforms are built for the special missions in many academies and aerospace institutions. Simulation platforms are not separated from models because model program is embedded into the platform. As models change, the platform inner structure program will be compiled again.
- (3) Most platforms run short of user interface. Only program structure about initialization, operation, and reprocessing is exposed by these platforms. It is not convenient for the experiment design or visual display.

3 BASIC ARCHITECTURE

Simulation platform is the key of overall simulation system. In order to meet the integration of different kinds of models, simulation platform should provide support for the bottom of different simulation applications. As it is said that, taking advantage of universalization, standardization and scalability of simulation platform, will reduce the difficulty of system development and the period of model establishment. Therefore, the quality of a platform is

the evaluation standard for a simulation system. For a simulation system, its universalization feature means that the platform is suited to many kinds of applications. In this paper, they are ground tests to missions. different kinds of aerospace Standardization emphasizes specificity among function-layers, criterion of connections between platform-platform and platform-model. Scalability is saying that the SIM simulator is prepared to cooperate with other simulation tools by performing different roles. There are two main scenarios: Integrating an external simulator within SIM or integrating the SIM framework within another simulator. Figure 1 shows the simulation system architecture.

SIM simulator consists of model management, experiment management, distribution management, operation management and node management.

Model management function can unify different kinds of models by making operations with the whole models or amending the information of a single model. Its design concentrates on layered-design theory. It realizes the separation between platform and models, model description and model realization. It is convenient for modeling, design of the experiment or the subsequent construction of simulation system.

Experiment management is mainly to assemble models with visualization software, configure connections among models, set model attributes. It establishes a communication bridge between user and simulation system. An excellent user interface does not only bring users enjoyable feelings, but also

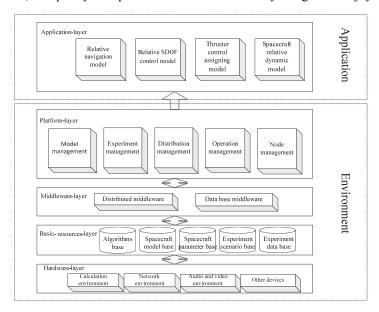


Figure 1: The simulation system architecture.

simplify the simulation experiment from foundation to testament. It is worth mentioning that layered-design theory is also used in the design concept of the experiment management. It separates models and model connections to increase flexibility of test design and reusability of models.

Distribution management adopts the idea of "distributed calculation, concentrated management". It distributes each model and model parameters to every node based on simulation scenarios. User can operate models which are on the node or the computer itself remotely. Related information will be shown on the test interface.

Simulation operation management is designed with responsibility for driving different kinds of models and attemperring distributedly. Meanwhile, it can also monitor the state of node and reserve data at breakpoint.

Node management is disposed on the calculation node computer with the function of node guard, model scheduling, and node state reporting. The function is in cooperation with manager node computer.

3.1 Model Management

Model management is the most basic and important function of platform. Each model works as a block. With the existing of model management function, simple blocks can be stacked to form the complex tall building and great mansion. Model management function drives and supervises models based on the specific characteristic.

3.1.1 Multipurpose Model-templet

In this paper, multipurpose model-templet is set up for the connections of different kinds of models. It doesn't only provide security for interaction between model and platform or model and model, but also provide convenience for modeling and model library construction in the future.

SIM can drive three kinds of models at present. We distinguish them by Model A, B and C.

Model A is a kind of simulation subsystem software which can be operating in Windows environment independently. It exists in the form of executable program(.exe). Connections of Model A include model basic information, model initial parameters, model input parameters, model output parameters and some files related to executable program. SIM drives Model A in the way of memory mapping. The file mapping to the memory which is marked by the name of model, can respond the control instruction to complete the function of model

initialization, step-by-step running, parameter modifying, stopping and so on.

The connection of Model B consists of three specialized functions. Data structure must be encapsulated by standard as the connection form which is satisfied with platform request. Model B exists by the form of dynamic link library (dll). This kind of model is so flexible that, it can be installed on any node. The form of model external port is "pInit, pInput, pOutput". Initialization is presented by a pointer "pInit". In the process of a simulation, some constant parameters which are used to describe model characteristics are pointed by "pInit". "pInput" is used to reserve parameters varied by time which are transferred to the model. By the same token, "pOutput" is used to reserve parameters the model exports. For example:

Initial Function:

Void XXX_Init (void *pInit, void *pInput, void *pOutput, void *pUser).

Steplike Function:

Void XXX_Sim (void *pInit, void *pInput, void *pOutput, void *pUser)

Reprocessing Function:

Void XXX_End (void *pInit, void *pInput, void *pOutput, void *pUser).

Where, XXX is the model name which can fully describe model function. "*pInit" is the pointer of model initial data. "*pInput/*pOutput" is the pointer describing the model input/output data. "*pUser" is the pointer to reserve the data which is defined by user himself. Three functions are set up to perform the model function of initialization, step-by-step running and stopping.

Model C presents the model in MATLAB. The connection consists of the same information as Model B.

3.1.2 Complex Model Driving

Generally speaking, a single variable can be reserved and assessed at any place by the computer. But structure variable is not like this. For different types of member variables, assessing process occurs at special place. Because every member variable is reserved in order in the memory.

The substance of data transmission between models is copying certain amount of memory space to another memory place. And model initialization parameters, input parameters and output parameters are encapsulated into structure. Each parameter is the member of the structure. In this case, at the beginning of data transferring, not only memory

space but also memory address of each parameter must to be known.

Different compiling environment, the detail of structure memory alignment is a little different, but has the same rule. Two concepts are introduced here.

The first one is variable alignment parameter. In Windows (32)/VC6.0 compiling environment, variable alignment parameter is just the variable accounting for the size in bytes. But in Linux compiling environment, some variable alignment parameters are based on operating system performance.

Except for variable alignment parameter, there is another alignment parameter called compiler alignment parameter (#pragma pack(n)). This numerical value can be set not only by code, but also modifying compiler property. In Windows (32)/VC6.0 compiling environment, "n" cannot be anything but 1/2/4/8, 8 is the default value. In Linux (32) GCC compiling environment, "n" only equals 1/2/4, its default value is 4.

After realizing the concepts of above-mentioned, two rules about structure memory alignment will be understood easily.

- 1) The character offset of each member in the structure to the first address must be integral multiple of the variable alignment parameter. If it is not, the memory will supply some related bytes after the last member.
- 2) The memory space occupied by the structure is integral multiple of alignment parameter. The same with last rule, if it does not meet the request, the memory will supply some related bytes after the last member. The alignment parameter in the sentence equals the less one of the biggest of all the members between the compiler alignment parameter.

Complex models are drived by the node model driving software. Based on each type of model-templet, SIM realizes the function of model initialization, operating, reprocessing, time hopping operation and cooperates experiment models with different steps assuring the coherence of model running.

Figure 2 shows the theory of Model A driving. Node model driving software writes data into the shared memory based on the protocol. Models read input data from the shared memory, and write output data into the shared memory after operating.

The connection of model B consists of three functions. Node model driving software drives models by the three functions. They are initial function, model operation control function and reprocessing function.

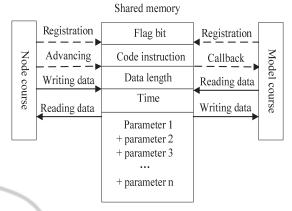


Figure 2: Node model driving software and theory of Model A driving.

To enhance the operating efficiency of simulation platform, large of initial jobs are done in the model initial function. 1) The platform can distribute memory space for all the model initial /input/output parameters. 2) The platform can find the memory address of each model input parameter in the structure according to the model connections. 3) The platform can reorder the model based on the model operation sequence. 4) The platform can drive each initial function of models.

Model operation control function determines the simulation cycle index based on platform simulation step and model step. Platform drives models orderly based on model operation sequence.

Reprocessing function invokes each reprocessing function of models, and releases the memory space opened before.

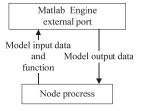


Figure 3: Schematic diagram of Model C driving.

Figure 3 is the schematic diagram of Model C driving. Node model driving software brings m(mdl) file to Matlab Engine to make calculation through Matlab Engine external interface, then reads computing results from Matlab Engine. In this mode, the relationship between node model driving software and Matlab Engine is C/S(Client/Service). Node model driving software regarding as a calculation servicer, through the distributed platform, sends or receives messages to Matlab Engine. The messages include orders and data information.

3.2 Other Function Management

Experiment management is to assemble models taking the form of visualization, configure model connections, distribute calculation nodes for models, model parameters, model simulation step, simulation condition, forming the simulation scenario.

User finishes configuring models in the form of graphical modeling, selecting models in the form of dragging, designing data stream in the form of connection. It becomes the concrete simulation experiment after assembling. Model is shown as a rectangular module with input/output ports. User finishes assembling models by the operation like clicking, dragging, and connecting. On the process of model assembling, if the input port doesn't match an output port, the platform will give an alarm. Models used to be assembled come from model library.

Distribution simulation management is designed using the thought of "distributed calculation, concentrated management". One of the main goals is distributing models and model parameters to each of related nodes, and operating models on the node or node computer itself in the long distance based on the simulation scenario. Simulation scenario is its input. After operating, user can see the operation tips on the user interface.

Operation management includes distributed scheduling, node supervising, simulation breakpoint reserving. It is the foundation of simulation system distributed operation. Operation management program is set in the management node computer. It can drive or distribute models on the calculation node. At the same time, it can supervise the node state in the long distance, and reserve the experiment information at breakpoint.

Node management is set on the calculation node. It is the foundation of "distributed calculation, concentrated management". Its main goals are node watching, driving the model on the node and reporting of the node state. It echoes operation management which is set on the management node. It drives and distributes models on the calculation node based on the simulation scenario file, and reports the state of local node to management node computer.

3.3 Modelling Distributed Environment

In the process of modeling for lunar orbit rendezvous mission, it is very important to consider the complicated spacecraft dynamic conditions, and high precision of the navigation and control in the close range rendezvous phase. These factors request sufficient accuracy of relates models. In the aspect of modeling, program developing must obey the standard of each type model port.

According to the requirements of the mission, we build relative navigation model, relative SDOF control model, thruster control assigning model, and spacecraft relative dynamic model.

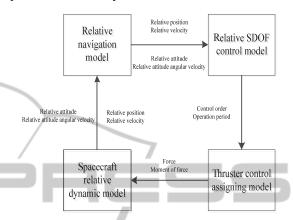


Figure 4: The information flow chart of models.

Figure 4 presents the information flow chart of models. At the beginning of the simulation, the first group input data is the initialization of the relative navigation model. After filtering calculation, measuring error is decreased, and the output data is the relative state between the spacecrafts. The relative state is as the input of the relative SDOF control model, through a series of control calculation, its output is the orders and operation period of the Thruster control assigning model thrusters. calculates force and moment of force based on the thruster orders. Spacecraft relative dynamic model inversely works out the relative state of spacecrafts which is controlled based on the output data from the thruster control assigning model. The relative state is regarded as the input on the next step.

In the process of modeling, we also use the model-templet to restrict the model ports. The above several models are developed in the form of dynamic link library. Model is named by capital English acronyms XXX. Each model includes the initial function XXX_init, Steplike function XXX_sim, and reprocessing function XXX_end. Three functions have the unified form. For example, the initial function of the relative navigation model is designed as Void RelNavigation_Init (void *pInit, void *pInput, void *pOutput, void *pUser). It is convenient for model management and model driving background.

Figure 5 is the class diagram of lunar orbit rendezvous simulation system.

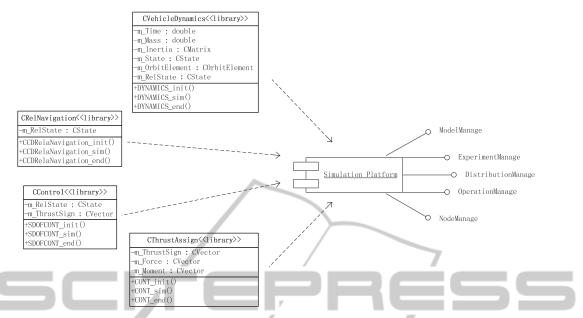


Figure 5: Class diagram of lunar orbit rendezvous simulation system.

From the class diagram, we can see that, the whole simulation system consists of spacecraft relative dynamic module, relative navigation module, relative SDOF control module, and thruster control assigning module. And, spacecraft relative dynamic module called CVehicleDynamics includes three functions which are DYNAMIC init, DYNAMIC sim and DYNAMIC end. Among them, the function of DYNAMIC sim includes many functions referring to the basic spacecraft dynamic. coordinate system transformation, orbital/attitude parameters transformation, spacecraft precision orbit determination. One of the main functions is the state transferring function which can transfer the absolute state of the spacecrafts to the relative state between them in the Hill coordinate. It is easy to analyze the problems of the spacecraft rendezvous at close range by the use of the relative state between two spacecrafts.

One of the main goals of relative navigation module which is named of CCDRelNavigation is decreasing the effects of measure noise to the docking sensor. It consists of three functions. They called CCDRelaNavigation init, CCDRelaNavigation sim and CCDRelaNavigation end. Moreover, CCDRelaNavigation sim function encapsulates some familiar filter methods. State transfer matrix and state equation mentioned in the filter method are worked out based on the spacecraft relative orbital/attitude equations.

Relative SDOF control module and thruster control assigning module form the chase control section. By the use of phase plane control method, CControl module receives navigation data, producing control signs and bring them to CThrustAssign module. CThrustAssign module gives out relative state data which is after the control.

As it is mentioned, SimulationPlatform presents SIM simulator. It has the function of model management, experiment management, distribution management, operation management and node management.

4 EXPERIMENTAL RESULTS

Making a test is a key link to verify the simulation platform. By analyzing the results of experiment, advantage and disadvantage of the platform and models will be found. It is beneficial for the platform to be developed and upgraded in the future. For models, we can also make some significant conclusions and found for modeling on the next step. At the same time, experiment results are the standard to estimate the validity of platform and models. In this section, we make a test about lunar orbit rendezvous simulation system.

The initialization of experiment is that target vehicle is moving in an approximately circular orbit around the moon. The orbit altitude is 300 kilometers. Its orbit angular velocity is 0.0007615

radian per second. The influence of the perturbative force to vehicles is not concerned, and attitude/orbit control motors maintain constant power. Initial

simulation parameters and initial model parameters are set as Table 1 and Table 2.

Table 1: Initial parameters of the simulation system.

Simulation initial condition	parameters
Starting time	0 s
System step	0.5 s
Step number	0
Terminal time	800 s
Current time	0 s
Model step	0.05 s

Table 2: Initial parameters of models.

Model initial condition	parameters
Relative position	[100, 0.01, 0.8] m
Relative velocity	[-0.155, -0.07, 0.01] m/s
Relative attitude angle	[2, 2, 2] deg
Relative attitude angle velocity	[0.02, 0.04, 0.02] deg/s
chase vehicle moment of inertia	[7285.46, 0.00, 0.00; 0.00, 6666.67, 0.00; 0.00, 0.00, 3285.47]
target vehicle moment of inertia	[7000.00, 0.00, 0.00; 0.00, 7000.00, 0.00; 0.00, 0.00, 5000.00]
chase vehicle mass	4000 kg
chase vehicle size	[1, 0, 0] m
Thruster force arm	[2.00, 2.00, 1.21] m
Position of the chase vehicle docking interface	[0.50, 2.00, 1.21] m
Position of the target vehicle docking interface	[0.50, 2.00, 1.21] m
Force of the orbit thruster	[25.00, 25.00, 25.00] N
Force of the attitude thruster	[10.00, 10.00, 10.00] N
Installation position of CCD camera	[1.00, 0.00, 0.00]
Viewing angle of CCD camera	[6, 6] deg
Viewing point of CCD camera	30 m
Position of mooring	30 m
Milestone of velocity	60 m
Mooring time	300 s
X position error(short/long range)	[0.02, 0.006; 0.15, 0.005] m
Y/Z position error(S/L range)	[0.02, 0.005; 0.2, 0.004] m
Vx error(S/L range)	[0.01, 0.006; 0.15, 0.01] m/s
Vy/Vz error(S/L range)	[0.01, 0.01; 0.2, 0.018] m/s
X attitude angle error(S/L range)	[0.23; 0.3] deg
Y/Z attitude angle error (S/L range)	[0.25; 0.3] deg
x attitude angle velocity (S/L range)	[0.17, 0.012; 0.2, 0.002] deg/s
y/z attitude angle velocity (S/L range)	[0.2, 0.02; 0.3, 0.005] deg/s

In view of the above test configuration, we make a test. The experiment results are as follow.

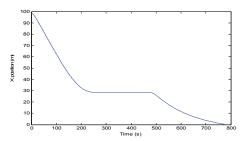


Figure 6: The X position transferring with time

- 1) Figure 6 shows the X position transferring with time. In this graph, we can see the relative position of the two spacecrafts changing from 100 meters to 0 meter in the end. In the distance of 30 meters, the chase vehicle stops for 300 seconds at predetermined location. The period of the whole docking process is about 780 seconds. Because the terminal time of simulation platform is 800 seconds, relative position is keeping 0 meter during the last 20 seconds.
- 2) Figure 7-Figure 8 respectively describes the process of relative Y/Z position changing with the docking distance. As approaching of two spacecrafts, the amplitude of the relative Y/Z position is decreasing, converging to 0 meter at last. In the process, target vehicle is staying in the CCD detecting scope of the chase vehicle all the time.

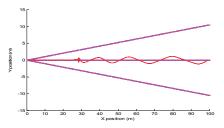


Figure 7: The Y position transferring with docking distance.

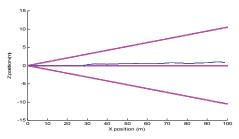


Figure 8: The Z position transferring with docking distance.

3) Figure 9-Figure 10 respectively describes the changing process of azimuth angle and pitch angle of CCD camera with the docking distance. The Attitude of Chase vehicle is controlled by impulses of attitude control engine continuously. When the distance between two spacecrafts is near, target vehicle is easy to deviate from the detection scope of the chase vehicle. By the correct flight strategy and suitable control mean, target vehicle doesn't deviate from the scope, but azimuth/pitch angle is going to converge in the range of ±0.5°. It illuminates the control strategy used in the model is available.

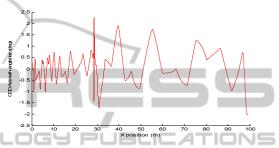


Figure 9: Azimuth angle of CCD camera.

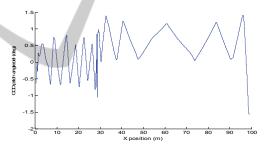


Figure 10: Pitch angle of CCD camera.

4) Figure 11-Figure 12 respectively shows the changing process of relative altitude angle and altitude angular velocity between two spacecrafts with the docking distance. At the end of the docking process, relative altitude angle precision can reach ±0.5 deg, relative altitude angular velocity precision can reach ±0.1 deg/s.

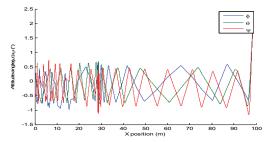


Figure 11: Relative altitude angle.

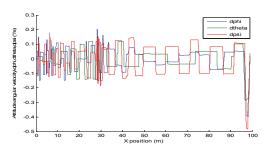


Figure 12: Altitude angular velocity.

5) Figure 13-Figure 15 respectively shows the changing process of velocity between two spacecrafts with the docking distance. At the end of the docking process, the relative velocity precision can be controlled in the range of $\pm 0.05 m/s$ which meets the requirements of spacecraft docking velocity.

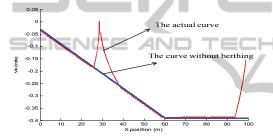


Figure 13: Changing process of Vx.

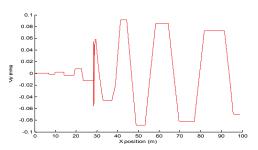


Figure 14: Changing process of Vy.

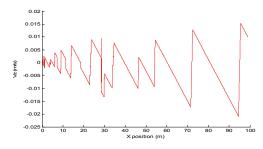


Figure 15: Changing process of Vz.

5 CONCLUSIONS

In this paper we present a modelling and simulation platform called SIM, which aims towards the study of parallel applications on distributed systems. This platform eases the process of designing and testing both the applications and the architectures.

The features of this platform are as follows. First, a great level of flexibility that allows modelling a wide range of designs is presented. Second, a friendly user interface, which helps to design the experiment and look for the architecture limits and bottlenecks is shown. Finally, the most important feature is a good separation between the platform and models which decreases much trouble in the process of system assembling.

The platform presents a modular design where the main components are the basic systems of distributed architecture such as computing, memory, storage, and network. This design also follows a hierarchical philosophy, where basic modules are grouped to compose bigger modules. SIM also provides several modules to simulate different components, and modeling strategies. Furthermore, the system allows the implementation of new modules by using a standardized interface.

Moreover, a series of models aon lunar orbit rendezvous has been performed. Relative navigation model, relative SDOF control model, thruster control assigning model, and spacecraft relative dynamic model have been configured and executed in SIM. Our platform shows very good result in the level of accuracy and performance obtained.

Future works on increasing the functionality of SIM are as following. First, a tool for modeling design and compiling. Second, encapsulation and decapsulation design for experiment models. Third, applying to multiple middleware protocols. Finally, accumulating and developing models. In the above case, we can lay the foundation of the manned lunar landing mission next step.

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