owl2proto: Enabling Semantic Processing in Modern Cloud Micro-Services

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Abstract: The usefulness of semantic technologies in the context of security has been demonstrated many times, e.g., for processing certification evidence, log files, and creating security policies. Integrating semantic technologies, like ontologies, in an automated workflow, however, is cumbersome since they introduce disruptions between the different technologies and data formats that are used. This is especially true for modern cloud-native applications, which rely heavily on technologies such as protobuf. In this paper we argue that these technology disruptions represent a major hindrance to the adoption of semantic technologies into the cloud and more effort and research is required to overcome them. We created one such approach called *owl2proto*, which provides an automatic translation of OWL ontologies into the protobuf data format. We showcase the seamless integration of an ontology and transmission of semantic data in an already existing cloud micro-service.

1 INTRODUCTION

Semantic technologies can establish a common understanding of, e.g., cloud concepts and their properties and thus have a high importance for the interoperability of cloud services. In the security context, semantic technologies have, for instance, been used to model certification evidence (Banse et al., 2021; Banse et al., 2023), to structure information in log files (Ben-Shimol et al., 2024), or to model general cloud security concepts (Takahashi et al., 2010).

We argue that while the academic discussion focuses on the semantic design of, e.g., cloud security concepts (Maroc and Zhang, 2019), the technological integration of semantic design and its technological implementation is lagging behind. Modern technologies like micro-services and RPCs, for example, are not integrated with technologies of the semantic web stack. We think that a better integration with such technologies would make semantic concepts easier to use, increase its adoption in different domains, and it could improve the interoperability of cloud systems, e.g. in multi-cloud and cloud-edge scenarios.

In this paper, we focus on *protobuf* as an example for this position and use the cloud security context as

an example application domain. Note, however, that our arguments apply beyond these examples. Protobuf is one of the most commonly used technologies for micro-services. Originally designed as a format to describe the serialization of network packages, it has evolved into an interface definition language, not only describing the exchanged data, but also the services that produce or consume this data.

Protobuf intentionally does not focus on the semantics of the exchanged data. Instead, it defines a syntax and structure of an object (called message), by describing which fields a programmer would use to fill this object. This includes primitive types, arrays and other messages. But the semantics of the data, such as, that it describes evidences gathered for a security incident, is beyond its scope. Developers would have to resort to storing this semantic information in other formats (e.g. RDF, JSON-LD, etc.) and then transmitting the actual data in a serialized form, creating a technology gap between the "semantic" world and the rest of the application.

In this paper, we argue for bridging the gap between semantic technologies and the integration with modern data processing. We demonstrate how to advance this integration by introducing a methodology and implementation that transforms ontology concepts into RPC definitions, and we point out use cases. The implemented tool is called *owl2proto*.

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2 BACKGROUND AND RELATED WORK

This section contains an overview of related work on semantic data processing, especially in the context of security and privacy. Additionally, we provide a background on the protobuf format our approach is based on.

2.1 Processing of Semantic Data

The term "semantic"—in the context of IT technology—refers to labeling data with common tags to enable their automatic recognition and processing. Ontologies provide one way to encode relationships between concepts, for example to represent hierarchies between them or other types of relationships. They can then be used by automatic technologies to put words, source code, and other objects into the right context. Several technologies exist to serialize such semantic data, e.g. to send them across communication channels. These include simple RDF serialization (Cyganiak et al., 2014), XML serialization of OWL (Hitzler et al., 2012) or more complex protocols such as JSON for Linking Data (JSON-LD) (Sporny et al., 2020).

In the context of security and privacy, semantic technologies have been used for malware analysis (Carvalho et al., 2016), authentication and authorization (Servos and Osborn, 2017), and in governance engineering (Esteves and Rodríguez-Doncel, 2024; Nadal et al., 2022). Even systems to automate the adherence to security compliance frameworks have been proposed (Banse et al., 2021; Banse et al., 2023).

2.2 gRPC and Protobuf

Protocol Buffers (or protobuf in short) is an opensource project to handle and transmit binary data in a structured way. It was originally developed internally in Google and released to the public in 2008. Protobuf's purpose is two-fold. First, it defines a binary wire format for the transmission of arbitrary data (so called *Messages*). It is very efficient and therefore used in scenarios where a high throughput of data is required. Second, protobuf can be seen as a *Interface Definition Language (IDL)*, which describes data structures and RPC interfaces in a (programming) language independent way.

These IDL files can then be used to auto-generate appropriate data structures in many chosen programming languages such as C/C++, Java, Python, Go and

others. Additionally, code for exposing the RPC interfaces as a service and consuming these as a client can also be auto-generated. Therefore, it is often used in the communication between a mesh of microservices, also in combination with additional RPC frameworks such as $gRPC¹$ or Connect², which take care of the actual network transmission of protobuf messages.

Listing 1: A protobuf message displaying some of the core concepts, such as messages, message options as well as the *oneof* keyword.

Listing 1 demonstrates several core concepts of the protobuf language as. First, every data structure in protobuf is called a *Message* (1). Each message must have a unique name (per package) and contains a list of fields. Each field also has a name, a type and a field number (the number after the equals sign). The type can either be one of the defined scalar types (such as string, $int32$ and others³) or another message (2) . The name as well as the field number must be unique inside their respective message. When protobuf messages are serialized and de-serialized, only field numbers are transmitted. Therefore, it is of utmost importance that field numbers are never changed or re-used for other fields, as long as compatibility needs to be ensured (e.g., within one major version).

To support this, the keyword reserved (4) can be used to denote that there was once a field with a

¹https://grpc.io

²https://connectrpc.com

³https://protobuf.dev/programming-guides/proto3/

Figure 1: Small excerpt of an ontology of cloud resources based on (Banse et al., 2023).

certain number (2 in this case) which is not used and cannot be used anymore. Protobuf also supports a special construct called oneof. It is somewhat similar to a *union* data structure in C. It defines a subgroup of fields (3) , in which only one of the defined fields (3a or 3b) can be set at a given time.

Finally, the notion of options can be seen as a system similar to annotations in other languages. They can provide additional meta-data to files, messages or fields. The contents of these are not transmitted over the wire, but they are available to both the transmitting and receiving party in the generated code. Next to a list of built-in options, also new options can be declared, as seen in \bullet . In this example, a new option note is declared which can be used to annotate a message $(1a)$.

3 MOTIVATION AND USE CASE

3.1 Use Case: Collecting Semantic Information for Cloud Security

To motivate the need for our approach, we considered existing previous works (see Section 2.1) from the security field that process semantic in the cloud. In the following, we will consider the design of (Banse et al., 2023) because they are already leveraging cloud technologies such as gRPC in addition to semantic data, but did not fully bridge the gap between both. They propose the collection of semantic information (also called *evidence*) about a cloud system, such as virtual resources and their properties. The aim is to assess certain properties of these assets with respect to security or privacy – in the context of certification. Because of the potential large amount of data gathered, their proposed system is split into different micro-services:

• a group of services (*discovery*) is responsible for collecting information about a cloud service and putting them into data objects described by an ontology (see Figure 1 for a small excerpt),

- a group of services (*evidence store*) is storing the gathered information in a common database,
- a group of services (*assessment*) is responsible for comparing the properties of such an evidence against a specific set of desired states using a rule engine.

All services use modern communication protocols, such as gRPC to communicate with each other, because of the high throughput needed. protobuf is also used to model and describe the service itself.

Listing 2: Excerpt of the assessment service API described in (Banse et al., 2023).

```
1 message Evidence {
2 string id = 1;
     google.protobuf.Struct evidence =←
        2;
\overline{4}5
6 message AssessEvidenceRequest {
     Evidence evidence = 1;
8 }
\overline{9}10 service Assessment {
11 rpc AssessEvidence (←
       AssessEvidenceRequest )
12 returns ( AssessEvidenceResponse );
13 }
```
Listing 2 shows a very small excerpt of this API definition, detailing the *Assessment* service, whose primary RPC call takes in an *Evidence*. While the evidence itself is described within the discovery services in terms of an ontology, the listing shows that this semantic information is partially lost. It is simply translated as a *Struct*, which is a protobuf definition of a generic key-value store. While this keeps the basic information – similar to a simple JSON – all other semantic aspects, such as entity inheritance or links to specific ontologies are lost. Instead a better approach would be dedicated messages for ontology types, like *Compute*, *VirtualMachine* and so on. Overall, it makes it very hard for users of this API because they need to have a look at the protobuf definition and additionally analyze how to model the contents of evidence/resource.

The only way to transmit such data and not loose semantic information would be to use formats like JSON-LD (Sporny et al., 2020) or other variants of JSON that are enriched with semantic information. This has several shortcomings (which are discussed in more detail in Section 5.2). Therefore, this service would benefit from a homogeneous solution in which the semantic data would also be available in the protobuf schema.

3.2 Generalization of the Use Case

Numerous similar cases can be found where differently named—but semantically similar—data need to be abstracted to the same level to enable, e.g., their aggregation. Consider the example of a social media platform for sports activities that allows users with different types of wearables to upload their activities including GPS coordinates, photos, their running statistics, etc. Aggregating data from different types of hardware and software manufacturers (e.g., smart watches, phones, specialized bike computers, etc.) implies the need for an ontological knowledge base that abstracts the most important properties and tools that support the work with this knowledge base in code. Other cases include, for instance, industrial data sharing platforms or log ingestion and analytics.

In general, the following requirements for designing a semantic data sharing approach for the cloud need to be considered:

- RQ1. Keeping semantic information, such as taxonomies and entity inheritance during the communication between different parts of the system (e.g. micro-services).
- RQ2. Seamless integration with cloud native technologies, such as gRPC/protobuf.
- RO3. Adaptability to the ever-changing world of DevOps and continuous integration and deployment.

4 OUR APPROACH: owl2proto

In trying to bridge the world of cloud native frameworks and ontology design, we choose protobuf as a means to convert ontology structures. Since protobuf is already a schema (and service) definition, it is very well suited to describe the exchanged data in a semantic way. We propose to auto-generate an appropriate protobuf schema out of a modelled ontology, specifically in the Web Ontology Language (OWL2) format. Therefore, we name our approach *owl2proto*.

Preparation Phase. Figure 2 shows the approach in detail. Two parties (*Sender* and *Receiver*) want to exchange semantically enriched data. The first step is to auto-generate a common protobuf definition out of the ontology that is shared by both parties $($ $\bullet)$). Execution of the translation process (see Section 4.1) is only necessary once and when the ontology changes (see Section 4.3).

Figure 2: Our approach on how to exchange semantic data between a sender and a receiver via a channel.

Semantic Data Exchange. Afterwards, sender $(2a)$ and receiver $(2b)$ can use the regular protobuf workflow to generate data structures in a programming language of their choice, e.g. Java/C++/Go or any other language supported by protobuf. Since protobuf allows the modular inclusion of different packages, it allows us to split the auto-generated parts and the manually modelled parts of the services into different source files. But in the end all protobuf files will be used in the code-generation of the respective server/client and exchange objects in the chosen programming language. The sender instantiates these data structures as memory objects and populates them with the RDF triplets (3) they want to send. After the sender initiates the communication $(4a)$, the receiver can de-serialize the transmitted values back into memory objects $(4b)$. Finally, the receiver can access the transmitted RDF triples (5) .

4.1 Translating Ontology Entities to Protobuf Constructs

One of the first steps is to map different concepts of OWL to protobuf. This includes mapping classes to messages and data properties as well as object properties to fields. We chose the OWL format (instead of RDF) because it already includes a structure containing classes and other concepts that are already more aligned with concepts in other programming languages. This eases the translation to protobuf, whereas a mapping of RDF entities would entail an additional layer of processing — which is essentially already done by OWL.

Listing 3: Excerpt of a generated protobuf output based on the ontology of Figure 1. A full version of the example is available in the GitHub repository.

```
1 option (owl.meta) = \{ 0 \}2 prefixes: [{
3 prefix: "ex"
4 iri: "http :// example.com/←-
      classes"
5 \, \}, {
6 prefix: "owl"
7 iri: "http :// www.w3.org/2002/07←-
       /owl#"
8 }]};
\overline{q}10 message VirtualMachine { 1
11 1a option (owl.class).iri = "ex:\leftrightarrowVirtualMachine ";
12 1b option (owl.class).parent = "ex: ←
      Compute";
13 option (owl.class).parent = "ex: ↔
      Resource";
14 option (owl.class).parent = "owl:←
       Thing";
15
16 2 string name = 1 [
17 2a (owl.property).iri = "ex:name",
18(2b) (owl.property).parent = "owl:\leftrightarrowtopDataProperty "
     19 2c (owl.property).class_iri = "ex:←-
      Resource"
                                    and the state of
20 \t 1;21
22 3 GeoLocation geo_location = 2 [
23 (owl.property).iri = "ex:has",
24 (owl.property).parent = "owl:←-
      topObjectProperty "
25 3c (owl.property).class_iri = "ex:\leftrightarrowCompute"
26 \t 1:
27
28 4 repeated string block_storage_ids←
       = 3 \text{ } 529 (owl.property).iri = "ex:\leftrightarrowhasMultiple ",
30 (owl.property).parent = "ex:has"
31 4c (owl.property).class_iri = "ex: ←
       VirtualMachine "
32 \t 1;
33 }
```
Listing 3 contains an excerpt of a generated protobuf output based on the ontology of Section 3.1. This illustrates the translation process as follows. For each OWL class, an associated protobuf message is created (1) for the VirtualMachine class).

Data Properties. For each data property used in the class, the (short) name of the property is taken as the name of the field. Since each property in OWL also has a type, we also need to map that accordingly. For primitive types, such as xsd:string, we perform a simple mapping to protobuf primitive types, such as string (2) .

Object Properties Referring to Blank Nodes. Properties which are only used in combination with blank (or anonymous) nodes, are set to the appropriate protobuf message type representing the OWL target class. This can be seen in 3. While the OWL object property is named has, we want to name the field according to the used OWL class GeoLocation (converted to geo location). Since GeoLocation is only used in blank nodes, we can directly use the translated protobuf message for the field's type.

Object Properties Referring to Identifiable Nodes. Properties that refer to other nodes which definitely have their own IRI and are identifiable are set to the string datatype and an id suffix is added to the protobuf field name. In this case, we only store the IRI referring to the other node instead of the content of the node itself in the final protobuf message. This can be seen in 4, where the hasMultiple object properties is converted to a string field of block storage ids since the target type in this class is BlockStorage⁴.

Retaining Semantic Structure Information. Protobuf message are well-suited to describe concepts like entities or OWL classes and their respective association to data / object properties. But, in order to retain the actual semantic information, such as IRIs, field types and other information contained in languages such as RDF or OWL, an additional step is needed. This entails the use of protobuf options in various forms. Listing 4 in the appendix contains the protobuf definition of these options. General information about the ontology are specified in file options, directly at the top of the protobuf file (0)). Ontology metadata about classes and properties are modelled as message annotations $(1a, 1b)$ and field annotations $(2a, 2b, 2c)$, respectively. Protobuf options are designed in a very efficient way and are not transmitted over the serialized channel. Instead the sending and receiving end can extract them out of fields and messages, as long as they keep their protobuf definitions in sync.

⁴Our Prototype implementation has an internal mapping which object properties are translated to a singular _{id} and which ones are translated to the plural *ids*. In the future, this could be taken from an annotation within the ontology.

4.2 Modelling Inheritance

Ontologies are inherently connected and use concepts like (multi)-inheritance. This means that a more concrete ontology entity will inherit the properties from all its parents. This concept does not exist in protobuf. Furthermore, some ontology entities (mainly non-leaf nodes) can be considered as interfaces in a programming language, since they hold denominators such as fields common to all their leaf nodes.

Therefore, we need to flatten the hierarchy of data properties in the individual ontology entities when translating them to protobuf messages. For example, if we look at our example ontology entity VirtualMachine, we can see that it derives from Compute and Resource. In this case, we need to include all data properties of both parent objects as protobuf fields, as shown in Listing 3 (\bigcirc , \bigcirc , 4). In order to keep the information which ancestor class specified the actual property, the option $(ow1.property)$.class_iri is used $(2c, 3c, 4c)$.

Furthermore, we want to keep the class hierarchy information. We make use of the fact that we can specify certain message options multiple times and include a (owl.class).parent option for each ancestor $(1b)$. For non-leaf nodes we make use of the oneof keyword to introduce a message that can be used similar to an interface.

4.3 Reacting to Changes

Every time the ontology changes, a new protobuf file needs to be generated using *owl2proto*. This step can (and should) be automated by a CI/CD system, such as GitHub actions. For example, a raw OWX file containing the OWL ontology could be stored alongside a code repository. Changes to this file can trigger a workflow that generates the protobuf files, so that developers can directly use it.

TECHNO

Since protobuf is a binary protocol, it uses field numbers to differentiate between different data fields within a message and not names. Therefore, to be compatible with previous versions of a protobuf file, field numbers must not change or be re-used. Otherwise, transmitted data will turn up in the wrong field. This has two implications:

- Multiple invocations of the translation process (on the same source file) need to produce the same output in the same order. Especially, the deduction of field numbers must be deterministic.
- If the source file changes, all fields that exist in both the original as well as the modified version of the file need to also have identical numbers.

There are several solutions to this problem, which we will discuss further, since each of them has their own specific drawback.

Non-Cryptographic Hash. One possible solution is to derive a unique field number for each field by its name, e.g. through the use of a cryptographic hash, such as $xxhash⁵$. But there are some caveats to consider:

- First, field numbers in protobuf are ranging from 1 to 536,870,911 (2^{29}) . The lowest xxhash implementation uses 32-bit, so we would need to further restrict the possible result space of hash 8 times, leading to more collisions.
- Certain field numbers (namely 19,000-19,999) are reserved for internal use
- Smaller numbers are more efficiently stored than larger numbers, contrasting a usual hash algorithm's way of using the whole result space to maximize entropy.

Read-in Previous Input. Another possibility is to read in the previous generated output as an input in order to assign existing fields the same number. For the initial generation, one could use a pre-defined ordering of fields based on lexical sorting of the field names and the parent's name which they are part of. Intentional space between groups of fields need to be left for extendability. If fields are removed, one can leverage the protobuf's reserved list in order to keep track of used field numbers. The greatest drawback of this solution is the need to have the previous version available, making it hard to maintain.

4.4 Implementation

As part of our research, we provide a prototype implementation of our approach as an Open-Source library⁶. The prototype is written in Go and largely uses the Go standard library to read in OWL in its XML variant (OWX) and generates the appropriate protobuf output.

With regards to the consistently problem (see Section 4.3), we decided to use a non-cryptographic hash function, such as xxhash in our prototype implementation, as this approach does not require a pre-existing proto file, which would have required some versioning functionality. To generate the unique field number, we use the name of the ontology class, the ontology property name and the names

⁵https://xxhash.com

⁶https://github.com/oxisto/owl2proto

of the parent classes as input for the hash function. The number of parent classes can be arbitrary. This ensures that the field numbers remain unique within the proto message, even if fields in different proto messages share the same name. An example is as follows. The input values for the field *name* (cf. Listing $3 \bullet$ and \bullet) are a tuple of (*VirtualMachine*,*Compute*,*Resource*,*name*).

Since the numerical range from 19000 to 19999 is reserved, it is imperative to map the resulting hash value to the range of 1 to 18999. This is accomplished by applying the mathematical modulo function. Although the numerical range from 20,000 to 536,870,911 is available for use, the range of 1 to 18,999 is adequate for the requirements of our ontology. Additionally, smaller numbers take up less space in protobuf serialization it is also more efficient.

5 DISCUSSION

In this section we discuss whether our approach can address the requirements that we elicited out of the use case presented in Section 3.1 and compare our approach to similar techniques, mainly JSON-LD.

5.1 Addressing the Use Case **Requirements**

RQ1. Keeping Semantic Information. We present an approach that can translate an ontology structure into a binary serialization format (protobuf), while keeping all the semantic information, even when transmitting the data. This is achieved through the combination of flattening hierarchies, translating certain types into identifiers and providing pseudointerfaces with oneof for certain entities. Finally, the use of message and field options allow us to embed all RDF/OWL concepts like IRIs and other annotations in the protobuf format⁷, making this requirement *fully addressed*.

RQ2. Seamless Integration. By choosing protobuf as our translation target, we aim at seamless integration into the cloud native world. We can easily integrate our generated messages into cloud native frameworks, such as gRPC or ConnectRPC since they use protobuf as their base. In contrast to other approaches like using gRPC for the general communication and JSON-LD for the serialization of seman-

tic data, we can offer one coherent API described in protobuf. This enables a more seamless developer experience and we consider this requirement *fully addressed*.

RQ3. Adaptability. Arguably this is the hardest requirement to fulfill. The pace of development in the modern cloud world can sometimes be astonishing. This also extends to the frequency of API and data model changes. Therefore, it also stretches to the development and enhancement of ontologies used in such services. On one side, our approach can easily be added into a CI/CD workflow, generating new protobuf files when ontology files change. On the other side, some inherit quirks of the protobuf format make auto-generation of protobuf files themselves quite hard. One such aspect is that protobuf message field numbers must not be changed once used. This makes it very hard for approaches that auto-generate those field numbers to a) consistently generate the same field numbers given the same input and b) to not re-assign already used field numbers if the input changes. While we present an initial approach using non-cryptographic hashes, our solution has several shortcomings that need further exploration. We therefore consider this requirement *partially addressed*.

5.2 Comparison with JSON-LD

When using systems that already rely on JSON as a data format, JSON-LD can provide a useful extension to increase interoperability. Regarding the technical usability, JSON-LD has some drawbacks. For example, JSON-LD files require dedicated libraries for processing. Furthermore, JSON-LD embeds or references the semantic data in each transmitted message which can significantly increase overhead. In comparison, protobuf is a binary format, which uses predefined structures making it more compact and faster to serialize/deserialize. Protobuf is therefore often considered to enable better performance.

Overall, JSON-LD can be considered an alternative to the protobuf-based approach presented here, but it presents significant drawbacks in comparison to protobuf. Protobuf, is the preferred choice for applications prioritizing performance and data integrity. Its binary format and pre-defined schema ensure efficient transmission and robust data consistency. Thus, it underlines the need for easy-to-use and efficient tooling that supports the integration of semantic data into modern (cloud) technologies.

 $⁷$ Although not tested, this could potentially even allow</sup> someone to re-construct a complete OWL file out of the generated protobuf file.

6 CONCLUSION AND FURTHER **RESEARCH**

In this paper we present *owl2proto*, an approach to bridge the world of ontology design with the world of modern cloud native frameworks and services. Our approach leverages protobuf, a popular format and framework for the description of data as well as a binary representation of the actual transfer. We show, using an example ontology based on cloud resources, that *owl2proto* addresses the requirements of seamlessly integrating semantic data into the Cloud.

We do note however, that there is further research to be conducted to fully address requirements of adaptability, stemming from complex requirements of protobuf itself. One possible approach could be to re-use techniques from the efforts to standardize RDF dataset canonicalization. Furthermore, projects like protovalidate⁸ could be used to translate OWL restrictions and constraints. Lastly, more validation of the approach itself using real-world ontologies and usecases needs to be performed.

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```
APPENDIX
```
Listing 4: The protobuf file defining our OWL options.

```
1 message EntityEntry {
2 string iri = 1;
     repeated string parent = 2;4 }
5 message PropertyEntry {
6 string iri = 1;
7 repeated string parent = 2;
8 string class_iri = 3;
9 }
10 message PrefixEntry {
11 string prefix = 1;
12 string iri = 2:
13 }
14 message Meta {
15 repeated PrefixEntry prefixes = 1;
16 }
17 extend google.protobuf. MessageOptions {
18 optional EntityEntry class = 50000;
19 }
20 extend google.protobuf. FieldOptions {
21 optional PropertyEntry property = 50000;
2223 extend google.protobuf. FileOptions {
24 optional Meta meta = 50000;
25 - 1
```
⁸https://github.com/bufbuild/protovalidate